

Regular Expressions

First, Some Announcements!

Second Midterm Logistics

- Our second midterm is next ***Monday, November 11th*** from ***7PM - 10PM***, locations TBA.
- Topic coverage is primarily lectures 06 – 13 (functions through induction) and PS3 – PS5. Finite automata and onward won't be tested here.
 - Because the material is cumulative, topics from PS1 – PS2 and Lectures 00 – 05 are also fair game.
- Room assignments to be posted late tonight.

Preparing for the Exam

- The top skills that will serve you well on this exam:
 - ***Knowing how to set up a proof***. This is a recurring theme across functions, sets, graphs, pigeonhole, and induction.
 - ***Distinguishing between assuming and proving***. This similarly cuts across all of these topics.
 - ***Reading new definitions***. This is at the heart of mathematical reasoning.
 - ***Writing proofs in line with definitions***. Folks often ask about whether they're being rigorous enough. Often "rigorous enough" simply means "following what the definitions say."
- Our personal recommendation: when working through practice problems, pay super extra close attention to these areas.

Preparing for the Exam

- As with the first midterm exam, we've posted a bunch of practice exams on the course website.
 - There are ten practice exams (yes, really!). We realistically don't expect anyone to complete them all. They're there to give you a feeling of what the exam might look like.
- Some general notes on preparing:
 - Q5 and Q6 on PS6, while technically on topics that aren't covered on the midterm, are great practice for the sorts of reasoning you'll need on the exam.
 - ***Keep the TAs in the loop when studying.*** Ask for feedback on any proofs you write when getting ready for the exam.
 - Don't skip on biological care and maintenance. Exams can be stressful, but please make time for basic things like showering, eating, etc. and for self-care in whatever form that takes for you.
- ***You can do this.*** Best of luck on the exam!

On to CS103!

Recap from Last Time

Regular Languages

- A language L is called a **regular language** if there is a DFA or an NFA for L .
- **Theorem:** The following are equivalent:
 - L is a regular language.
 - There is a DFA D where $\mathcal{L}(D) = L$.
 - There is an NFA N where $\mathcal{L}(N) = L$.
- In other words, knowing any one of the above three facts means you know the other two.

Language Concatenation

- If $w \in \Sigma^*$ and $x \in \Sigma^*$, then wx is the **concatenation** of w and x .
- If L_1 and L_2 are languages over Σ , the **concatenation** of L_1 and L_2 is the language L_1L_2 defined as

$$L_1L_2 = \{ x \mid \exists w_1 \in L_1. \exists w_2 \in L_2. x = w_1w_2 \}$$

- Example: if $L_1 = \{ a, ba, bb \}$ and $L_2 = \{ aa, bb \}$, then

$$L_1L_2 = \{ aaa, abb, baaa, babb, bbaa, bbbb \}$$

Lots and Lots of Concatenation

- Consider the language $L = \{ \mathbf{aa}, \mathbf{b} \}$
- LL is the set of strings formed by concatenating pairs of strings in L .

$\{ \mathbf{aaaa}, \mathbf{aab}, \mathbf{baa}, \mathbf{bb} \}$

- LLL is the set of strings formed by concatenating triples of strings in L .

$\{ \mathbf{aaaaaa}, \mathbf{aaaab}, \mathbf{aabaa}, \mathbf{aabb}, \mathbf{baaaa}, \mathbf{baab}, \mathbf{bbaa}, \mathbf{bbb} \}$

- $LLLL$ is the set of strings formed by concatenating quadruples of strings in L .

$\{ \mathbf{aaaaaaaa}, \mathbf{aaaaaab}, \mathbf{aaaabaa}, \mathbf{aaaabb}, \mathbf{aabaaaa}, \mathbf{aabaab}, \mathbf{aabbaa}, \mathbf{aabbb}, \mathbf{baaaaaa}, \mathbf{baaaab}, \mathbf{baabaa}, \mathbf{baabb}, \mathbf{bbaaaa}, \mathbf{bbaab}, \mathbf{bbbaa}, \mathbf{bbbb} \}$

Language Exponentiation

- We can define what it means to “exponentiate” a language as follows:

$$L^0 = \{\varepsilon\} \quad L^{n+1} = LL^n$$

- So, for example, $\{ \mathbf{aa}, \mathbf{b} \}^3$ is the language

$$\{ \mathbf{aaaaaa}, \mathbf{aaaab}, \mathbf{aabaa}, \mathbf{aabb}, \\ \mathbf{baaaa}, \mathbf{baab}, \mathbf{bbaa}, \mathbf{bbb} \}$$

The Kleene Closure

- An important operation on languages is the ***Kleene Closure***, which is defined as

$$L^* = \{ w \in \Sigma^* \mid \exists n \in \mathbb{N}. w \in L^n \}$$

- Mathematically:

$$w \in L^* \quad \text{iff} \quad \exists n \in \mathbb{N}. w \in L^n$$

- Intuitively, all possible ways of concatenating zero or more strings in L together, possibly with repetition.

The Kleene Closure

If $L = \{ \mathbf{a}, \mathbf{bb} \}$, then $L^* = \{$

$\epsilon,$

$\mathbf{a}, \mathbf{bb},$

$\mathbf{aa}, \mathbf{abb}, \mathbf{bba}, \mathbf{bbbb},$

$\mathbf{aaa}, \mathbf{aabb}, \mathbf{abba}, \mathbf{abbbb}, \mathbf{bbaa}, \mathbf{bbabb}, \mathbf{bbbba}, \mathbf{bbbbbb},$

\dots

$\}$

Think of L^* as the set of strings you can make if you have a collection of rubber stamps – one for each string in L – and you form every possible string that can be made from those stamps.

Closure Properties

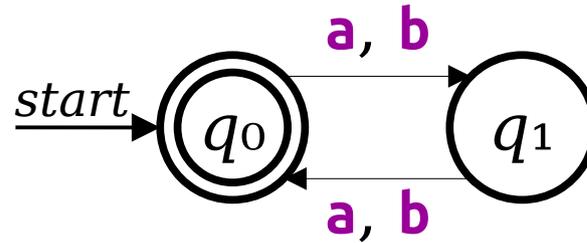
- ***Theorem:*** If L_1 and L_2 are regular languages over an alphabet Σ , then so are the following languages:
 - $L_1 \cup L_2$
 - L_1L_2
 - L_1^*
- These (and other) properties are called ***closure properties of the regular languages.***

New Stuff!

Another View of Regular Languages

Devices for Articulating Regular Languages

- Finite Automata



- Set (or other Mathematical) Notation

$\{ w \in \Sigma^* \mid w\text{'s length is even} \}$

- State Transition Table

	a	b
q_0	q_1	q_1
q_1	q_0	q_0

- ***New!*** Regular Expressions

Devices for Articulating Regular Languages

- Finite Automata



- Set (or other Mathematical) Notation

$\{ w \in \Sigma^* \mid w\text{'s length is even} \}$

- State Transition

Note: This one is not unique to regular languages! We can express non-regular languages with set builder notation, as well. More on that another day, when we explore other families of languages.

- **New!** Regular

Regular Expressions

- ***Regular expressions*** are a way of describing a language via a string representation.
- They're used just about everywhere:
 - They're built into the JavaScript language and used for data validation.
 - They're used in the UNIX grep and flex tools to search files and build compilers.
 - They're employed to clean and scrape data for large-scale analysis projects.
- Conceptually, regular expressions are strings describing how to assemble a larger language out of smaller pieces.

Rethinking Regular Languages

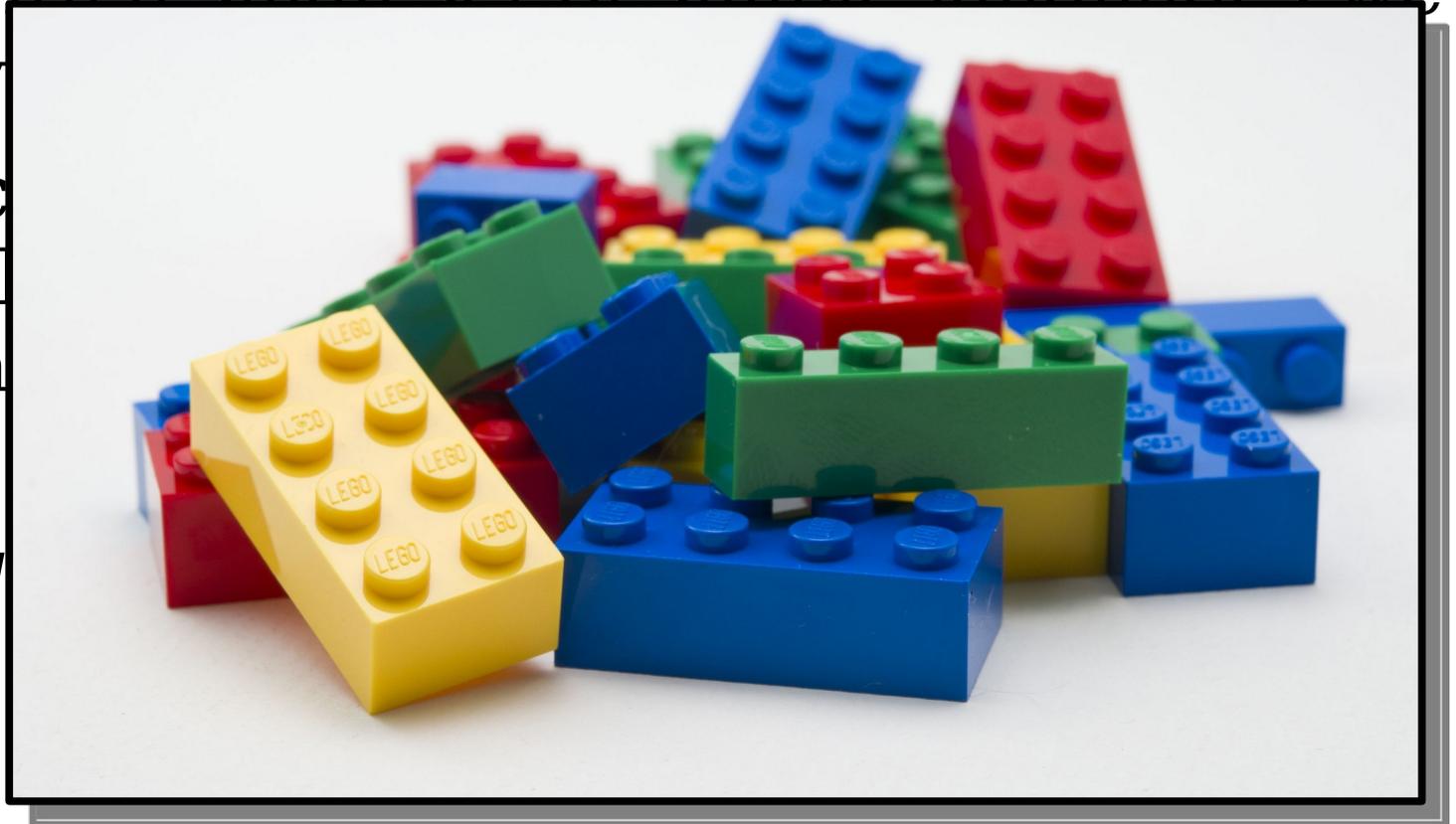
- We currently have several tools for showing a language L is regular:
 - Construct a DFA for L .
 - Construct an NFA for L .
 - Combine several simpler regular languages together via closure properties to form L .
- We have not spoken much of this last idea.

Constructing Regular Languages

- **Idea:** Build up all regular languages as follows:
 - Start with a small set of simple languages we already know to be regular.
 - Using closure properties, combine these simple languages together to form more elaborate languages.
- *This is a bottom-up approach to the regular languages.*

Constructing Regular Languages

- **Idea:** Build up all regular languages as follows:
 - Start with a small set of simple languages we already know
 - Using closure operations (union, concatenation, Kleene star) to elaborate on these
- *This is a regular language*



Atomic Regular Expressions

- The regular expressions begin with three simple building blocks.
- The symbol \emptyset is a regular expression that represents the empty language \emptyset .
- For any $a \in \Sigma$, the symbol a is a regular expression for the language $\{a\}$.
- The symbol ϵ is a regular expression that represents the language $\{\epsilon\}$.
 - **Remember:** $\{\epsilon\} \neq \emptyset!$
 - **Remember:** $\{\epsilon\} \neq \epsilon!$

Compound Regular Expressions

- If R_1 and R_2 are regular expressions, R_1R_2 is a regular expression for the *concatenation* of the languages of R_1 and R_2 .
- If R_1 and R_2 are regular expressions, $R_1 \cup R_2$ is a regular expression for the *union* of the languages of R_1 and R_2 .
- If R is a regular expression, R^* is a regular expression for the *Kleene closure* of the language of R .
- If R is a regular expression, (R) is a regular expression with the same meaning as R .

Operator Precedence

- Here's the operator precedence for regular expressions:

(R)

R^*

R_1R_2

$R_1 \cup R_2$

- So **ab*cUd** is parsed as **((a(b*))c)Ud**

Regular Expression Examples

- The regular expression **trickUtreat** represents the language

{ **trick**, **treat** }.

- The regular expression **boo*** represents the regular language

{ **boo**, **booo**, **boooo**, ... }.

- The regular expression **candy!(candy!)*** represents the regular language

{ **candy!**, **candy!candy!**, **candy!candy!candy!**, ... }.

Regular Expressions, Formally

- The *language of a regular expression* is the language described by that regular expression.
- Formally:
 - $\mathcal{L}(\epsilon) = \{\epsilon\}$
 - $\mathcal{L}(\emptyset) = \emptyset$
 - $\mathcal{L}(a) = \{a\}$
 - $\mathcal{L}(R_1R_2) = \mathcal{L}(R_1) \mathcal{L}(R_2)$
 - $\mathcal{L}(R_1 \cup R_2) = \mathcal{L}(R_1) \cup \mathcal{L}(R_2)$
 - $\mathcal{L}(R^*) = \mathcal{L}(R)^*$
 - $\mathcal{L}((R)) = \mathcal{L}(R)$

Worthwhile activity: Apply this recursive definition to

$a(b \cup c)((d))$

and see what you get.

Designing Regular Expressions

- Let $\Sigma = \{a, b\}$.
- Let $L = \{ w \in \Sigma^* \mid w \text{ contains } aa \text{ as a substring} \}$.

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bbabbbaabab

aaaa

bbbbbabbbbaabbbb

Designing Regular Expressions

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bbabbb**aa**bab

aaa

bbbbbabbbb**aa**bbbbbb

Designing Regular Expressions

- Let $\Sigma = \{a, b\}$.
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$\Sigma^*aa\Sigma^*$

bbabbb**aa**bab

aaa

bbbbbabbbb**aa**bbbbbb

Designing Regular Expressions

- Let $\Sigma = \{a, b\}$.
- Let $L = \{ w \in \Sigma^* \mid |w| = 4 \}$.

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The length of
a string w is
denoted $|w|$

Designing Regular Expressions

- Let $\Sigma = \{a, b\}$.
- Let $L = \{ w \in \Sigma^* \mid |w| = 4 \}$.

Write a regex for this language.

Answer at

<https://cs103.stanford.edu/pollev>

Designing Regular Expressions

- Let $\Sigma = \{a, b\}$.
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$\Sigma\Sigma\Sigma\Sigma$

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$\Sigma\Sigma\Sigma\Sigma$

aaaa
baba
bbbb
baaa

Designing Regular Expressions

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$\Sigma\Sigma\Sigma\Sigma$

aaaa

baba

bbbb

baaa

Designing Regular Expressions

- Let $\Sigma = \{a, b\}$.
- Let $L = \{ w \in \Sigma^* \mid |w| = 4 \}$.

Σ^4

aaaa
baba
bbbb
baaa

Designing Regular Expressions

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Σ^4

aaaa
baba
bbbb
baaa

Designing Regular Expressions

- Let $\Sigma = \{a, b\}$.
- Let $L = \{ w \in \Sigma^* \mid w \text{ contains at most one } a \}$.

Here are some candidate regular expressions for the language L . Which of these are correct?

$\Sigma^*a\Sigma^*$

$b^*ab^* \cup b^*$

$b^*(a \cup \epsilon)b^*$

$b^*a^*b^* \cup b^*$

$b^*(a^* \cup \epsilon)b^*$

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Designing Regular Expressions

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$b^*(a \cup \epsilon)b^*$

bbbbabbb

bbbbbb

abbb

a

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$b^*(a \cup \epsilon)b^*$

bbbbabbb

bbbbbb

abbb

a

Designing Regular Expressions

- Let $\Sigma = \{a, b\}$.
- Let $L = \{ w \in \Sigma^* \mid w \text{ contains at most one } a \}$.

$b^*a?b^*$

bbbbabbb

bbbbbb

abbb

a

A More Elaborate Design

- Let $\Sigma = \{ \mathbf{a}, ., @ \}$, where \mathbf{a} represents “some letter.”
- Let's make a regex for email addresses.

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- Let $\Sigma = \{ \mathbf{a}, \mathbf{.}, \mathbf{@} \}$, where \mathbf{a} represents “some letter.”
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cs103@cs.stanford.edu
first.middle.last@mail.site.org
dot.at@dot.com

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aa*

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aa* (.aa*)*

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$\mathbf{aa^* (.aa^*)^*}$

$\mathbf{cs103@cs.stanford.edu}$

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A More Elaborate Design

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aa* (**.aa***)* **@**

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aa* (**.aa***)* **@** **aa*.aa***

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aa* (**.aa***)* **@** **aa*.aa*** (**.aa***)*

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aa* (**.aa***)* @ **aa*.aa*** (**.aa***)*

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A More Elaborate Design

- Let $\Sigma = \{ a, ., @ \}$, where **a** represents “some letter.”
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a⁺ (**.**aa*)* @ aa* .aa* (**.**aa*)*

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first.middle.last@mail.site.org
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A More Elaborate Design

- Let $\Sigma = \{ a, ., @ \}$, where **a** represents “some letter.”
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a⁺ **(.aa*)*** **@** **aa*.aa*** **(.aa*)***

cs103@cs.stanford.edu
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A More Elaborate Design

- Let $\Sigma = \{ \mathbf{a}, \mathbf{.}, \mathbf{@} \}$, where \mathbf{a} represents “some letter.”
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$\mathbf{a}^+ \mathbf{(.a^+)^* @ a^+ .a^+ (.a^+)^*}$

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 $\mathbf{first.middle.last@mail.site.org}$
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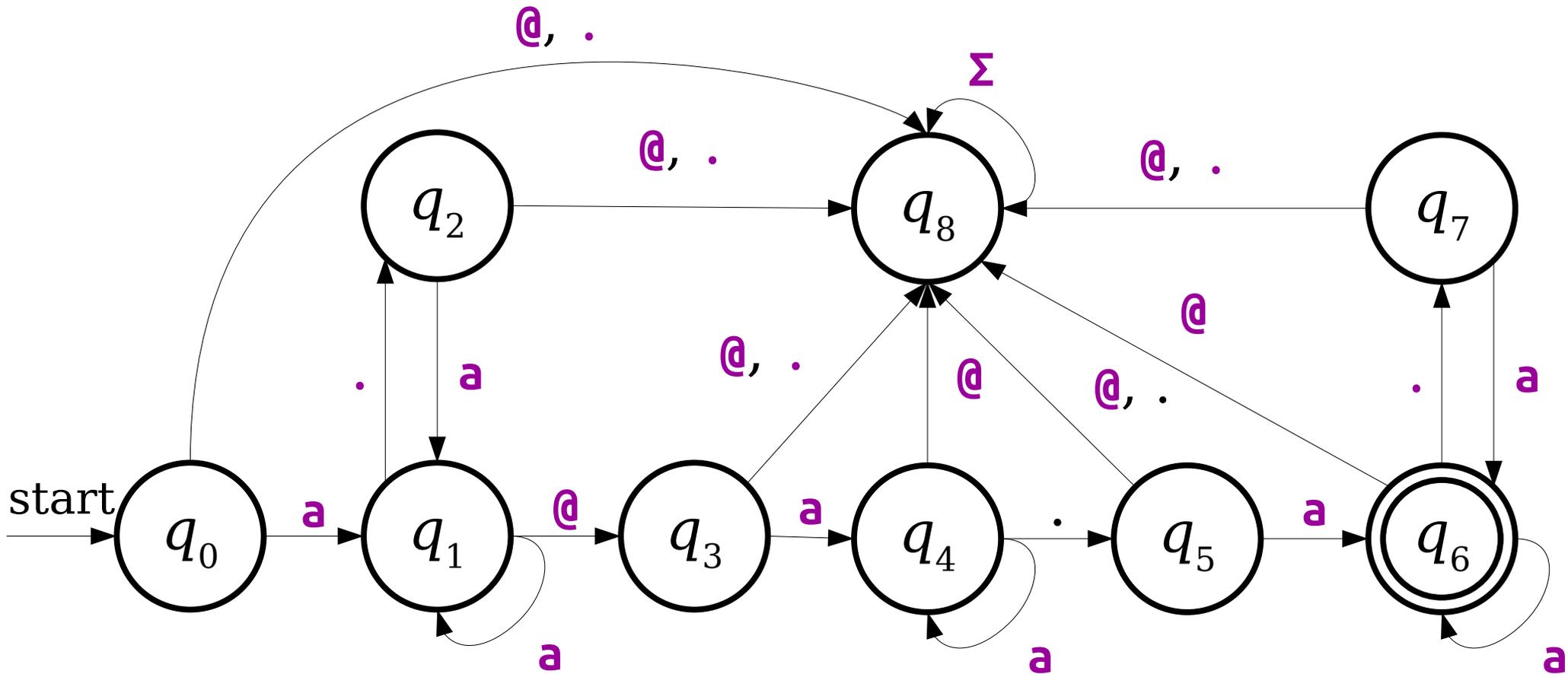
- Let $\Sigma = \{ a, ., @ \}$, where **a** represents “some letter.”
- Let's make a regex for email addresses.

a⁺ (**.****a**⁺)^{*} **@** **a**⁺ (**.****a**⁺)⁺

cs103**@****cs**.stanford.edu
first.**middle**.**last****@**mail.site.org
dot.**at****@**dot.com

For Comparison

$a^+ (.a^+) * @a^+ (.a^+)^+$



Shorthand Summary

- R^n is shorthand for $RR \dots R$ (n times).
 - Edge case: define $R^0 = \varepsilon$.
- Σ is shorthand for “any character in Σ .”
- $R?$ is shorthand for $(R \cup \varepsilon)$, meaning “zero or one copies of R .”
- R^+ is shorthand for RR^* , meaning “one or more copies of R .”

The Lay of the Land

Languages you can
build a DFA for.

Languages you can
build an NFA for.

***Regular
Languages***

```
graph TD; A[Languages you can build a DFA for.] --> C((Regular Languages)); B[Languages you can build an NFA for.] --> C;
```

Languages you can
build a DFA for.

Languages you can
build an NFA for.

Regular Languages

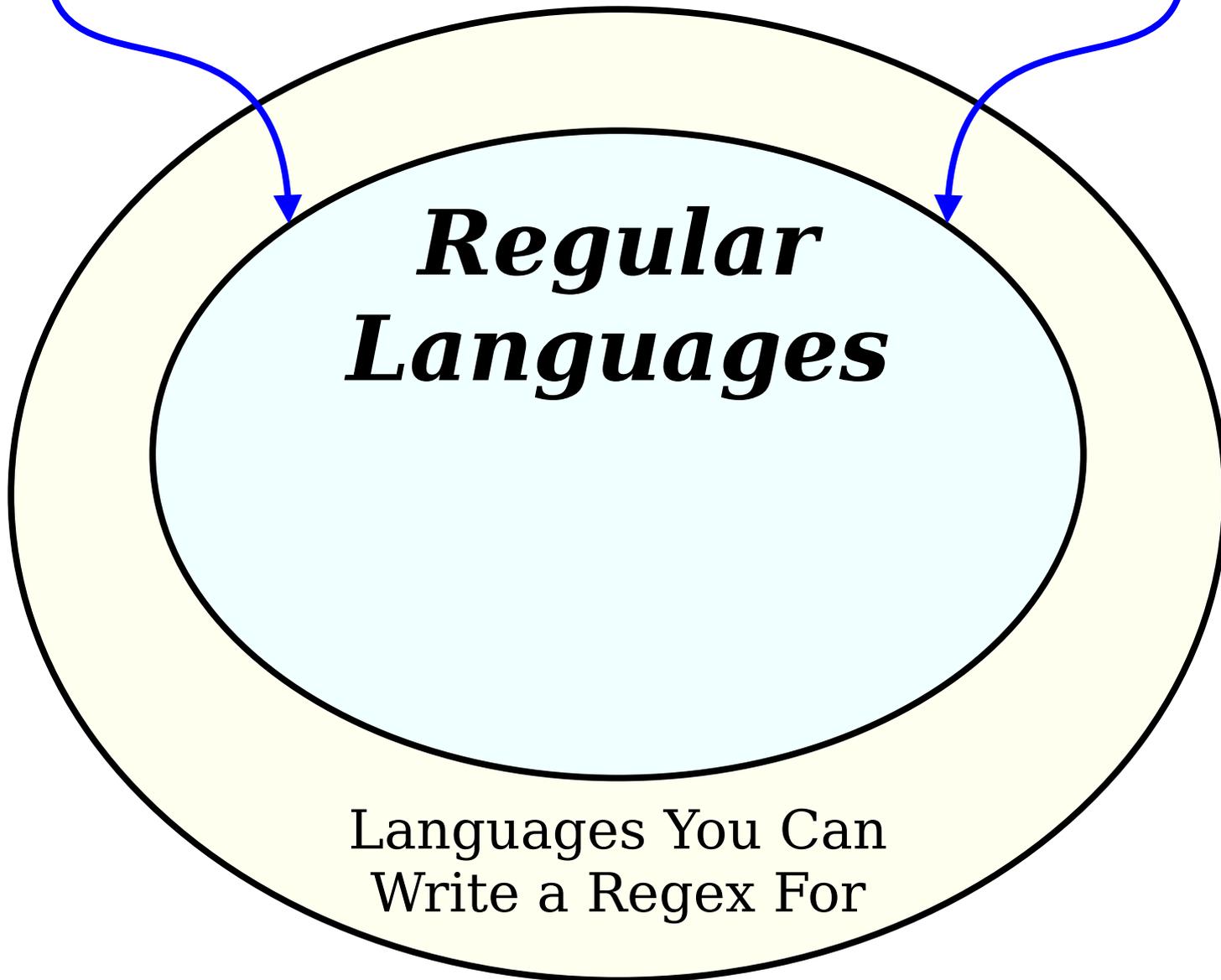
Languages You Can
Write a Regex For

Languages you can
build a DFA for.

Languages you can
build an NFA for.

***Regular
Languages***

Languages You Can
Write a Regex For



Languages you can
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***Regular
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Regular Languages

Languages You Can
Write a Regex For

The Power of Regular Expressions

Theorem: If R is a regular expression, then $\mathcal{L}(R)$ is regular.

Proof idea: Use induction!

- The atomic regular expressions all represent regular languages.
- The combination steps represent closure properties.
- So anything you can make from them must be regular!

Thompson's Algorithm

- In practice, many regex matchers use an algorithm called ***Thompson's algorithm*** to convert regular expressions into NFAs (and, from there, to DFAs).
 - Read Sipser if you're curious!
- ***Fun fact:*** the “Thompson” here is Ken Thompson, one of the co-inventors of Unix!

Languages you can
build a DFA for.

Languages you can
build an NFA for.

***Regular
Languages***

```
graph TD; A[Languages you can build a DFA for.] --> C((Regular Languages)); B[Languages you can build an NFA for.] --> C;
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Languages you can
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Regular Languages

Languages You Can
Write a Regex For

Languages you can
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Regular Languages

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The Power of Regular Expressions

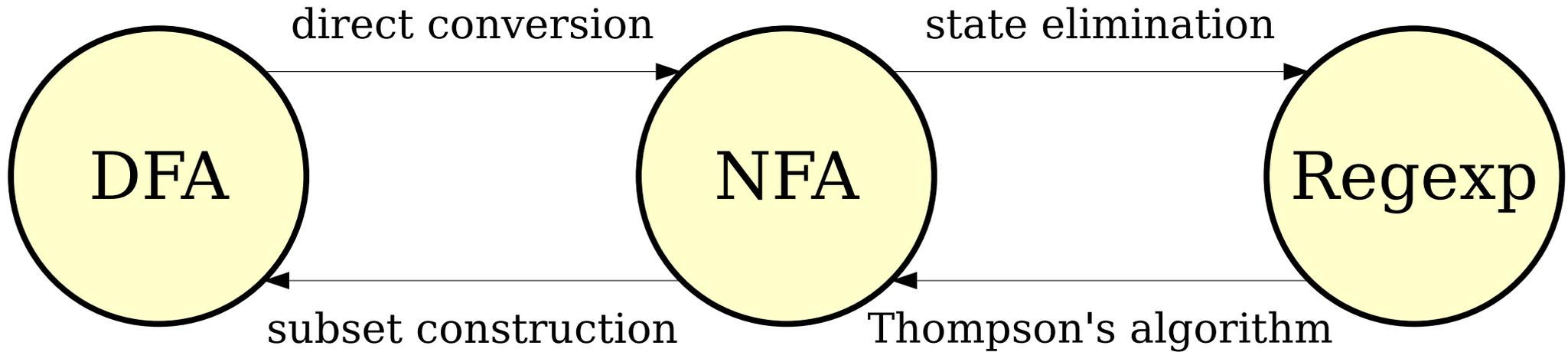
Theorem: If L is a regular language, then there is a regular expression for L .

This is not obvious!

Proof idea: Show how to convert an arbitrary NFA into a regular expression.

See appendix!
(“state elimination”)

Our Transformations



Theorem: The following are all equivalent:

- L is a regular language.
- There is a DFA D such that $\mathcal{L}(D) = L$.
- There is an NFA N such that $\mathcal{L}(N) = L$.
- There is a regular expression R such that $\mathcal{L}(R) = L$.

Why This Matters

- The equivalence of regular expressions and finite automata has practical relevance.
 - Regular expression matchers have all the power available to them of DFAs and NFAs.
- This also is hugely theoretically significant: the regular languages can be assembled “from scratch” using a small number of operations!

Your Action Items

- ***Read “Guide to Regexes”***
 - There’s a lot of information and advice there about how to write regular expressions, plus a bunch of worked exercises.
- ***Read “Guide to State Elimination”***
 - It’s a beautiful algorithm. The Guide goes into a lot more detail than what we did here.

Next Time

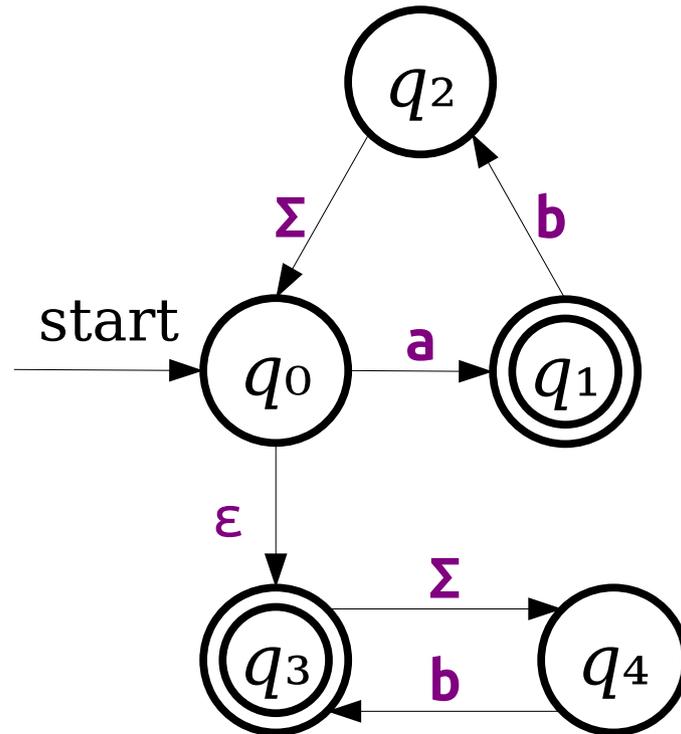
- ***Intuiting Regular Languages***
 - What makes a language regular?
- ***The Myhill-Nerode Theorem***
 - The limits of regular languages.

Appendix:

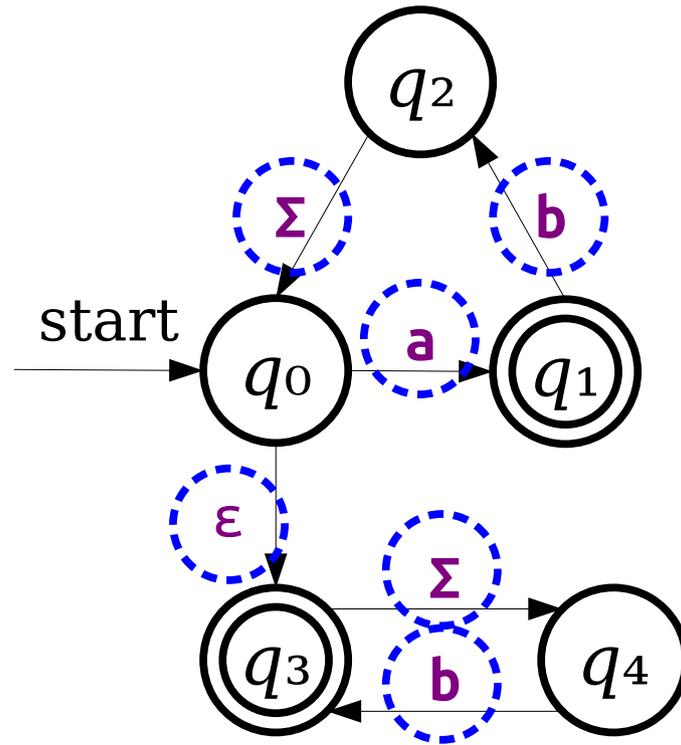
State Elimination (NFA \rightarrow Regex)

See also: *Guide to State Elimination*

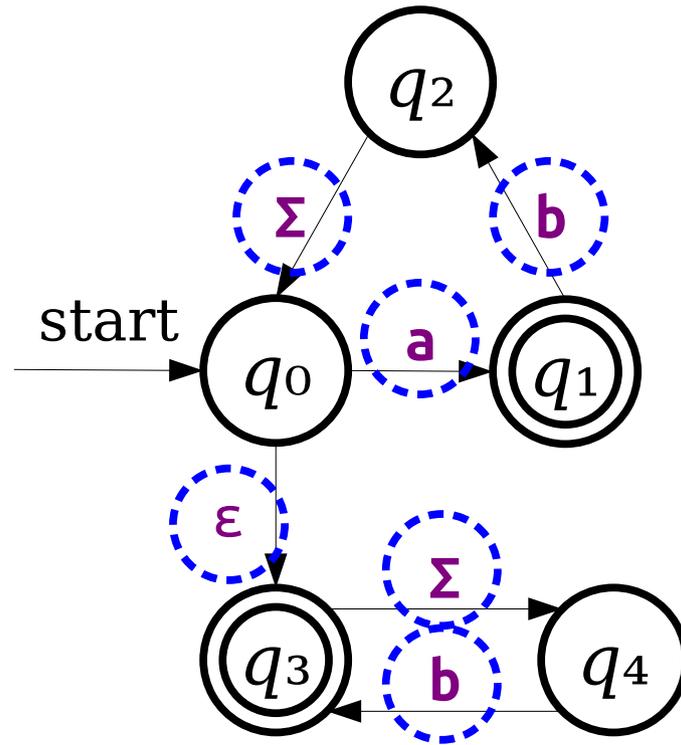
Generalizing NFAs



Generalizing NFAs

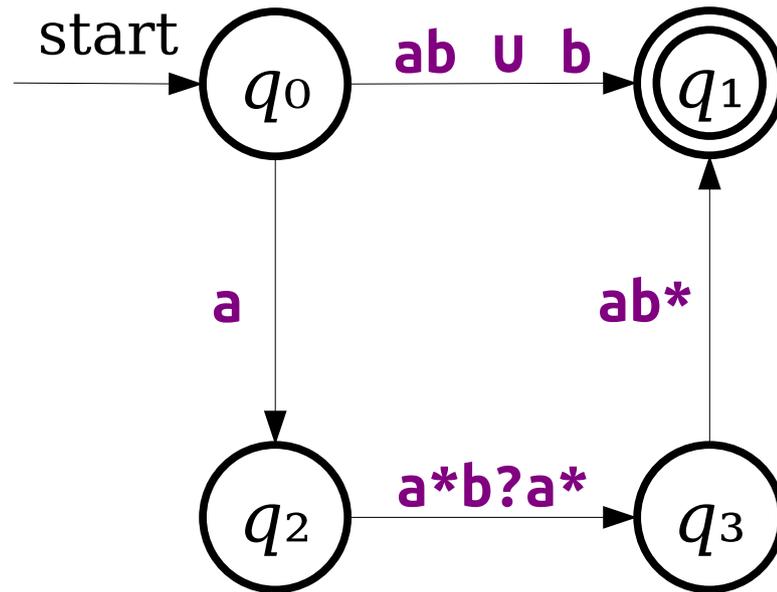


Generalizing NFAs

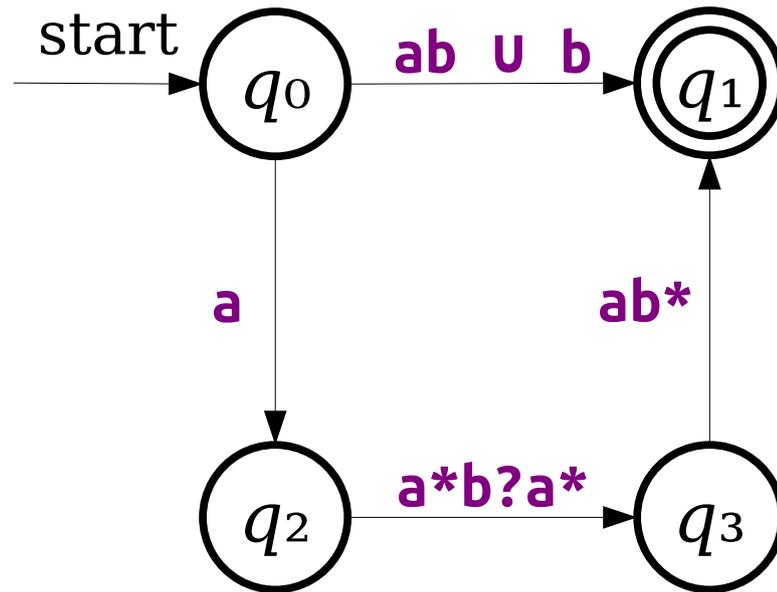


These are all regular expressions!

Generalizing NFAs

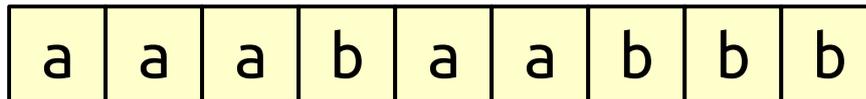
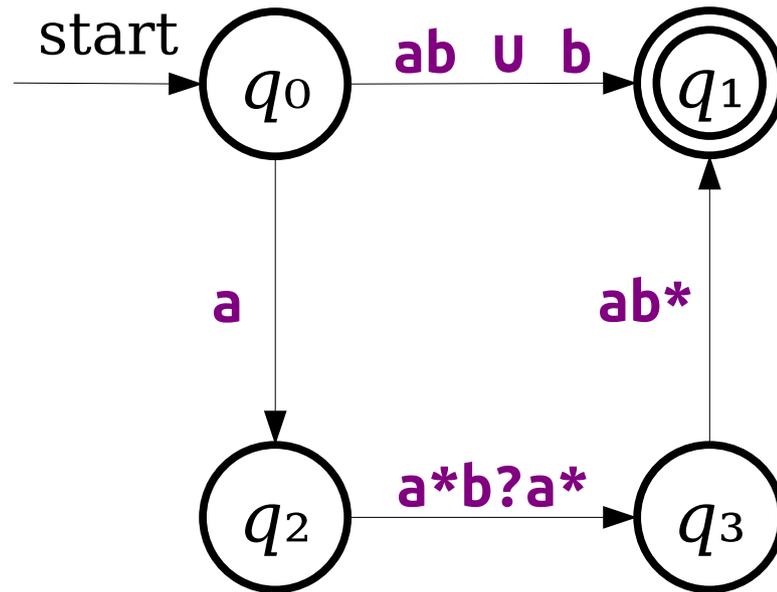


Generalizing NFAs

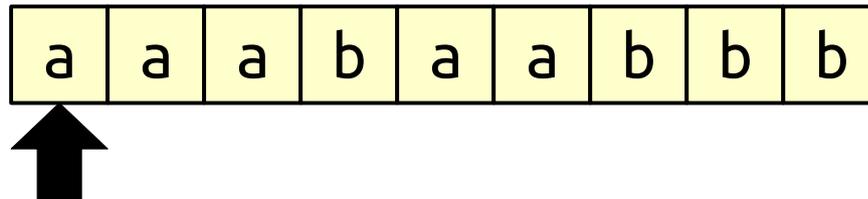
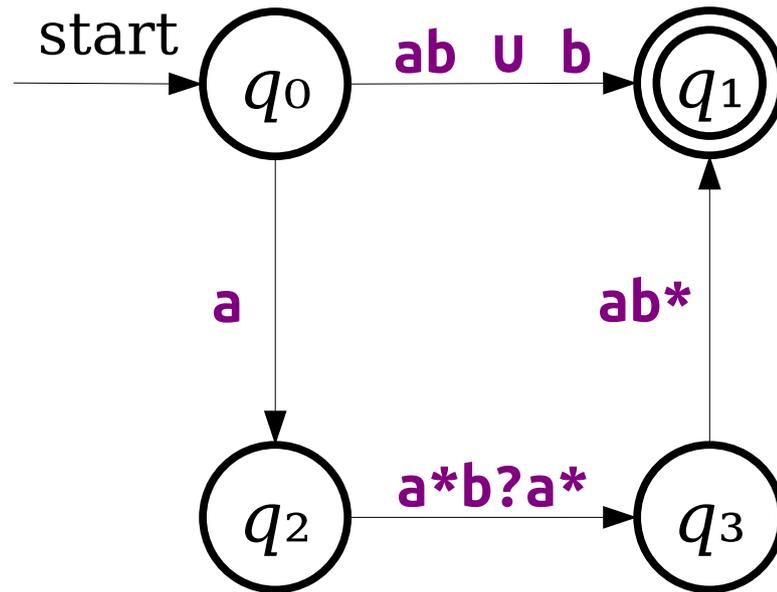


Note: Actual NFAs aren't allowed to have transitions like these. This is just a thought experiment.

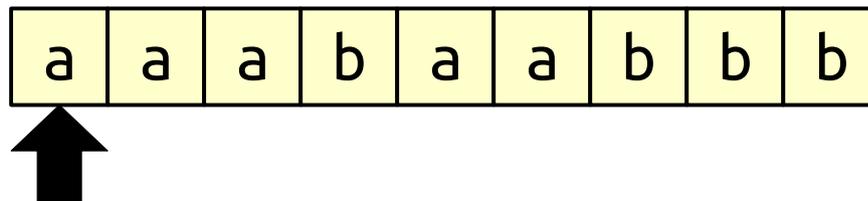
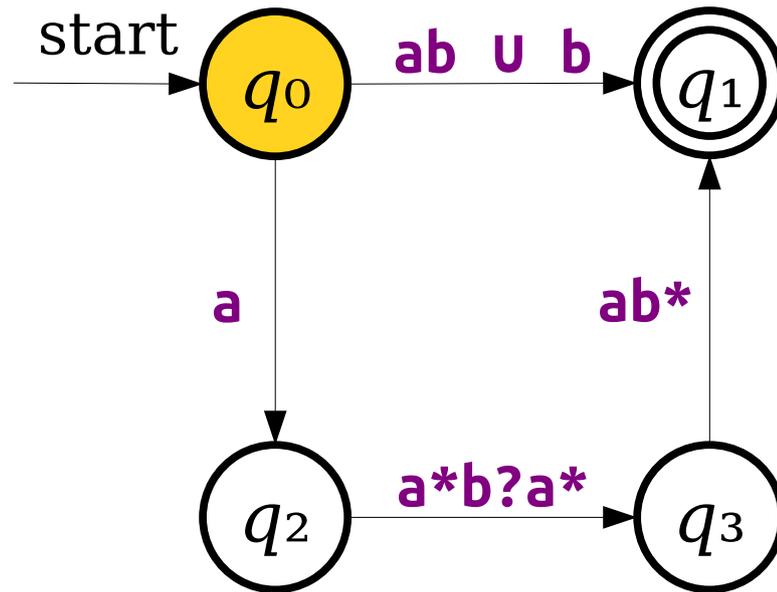
Generalizing NFAs



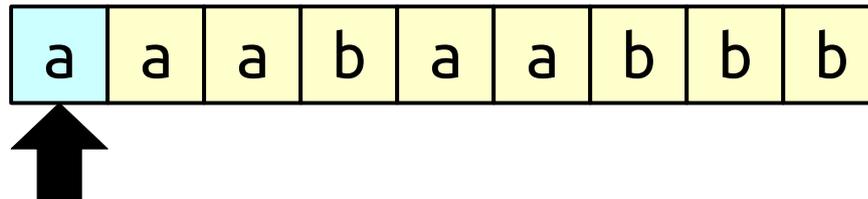
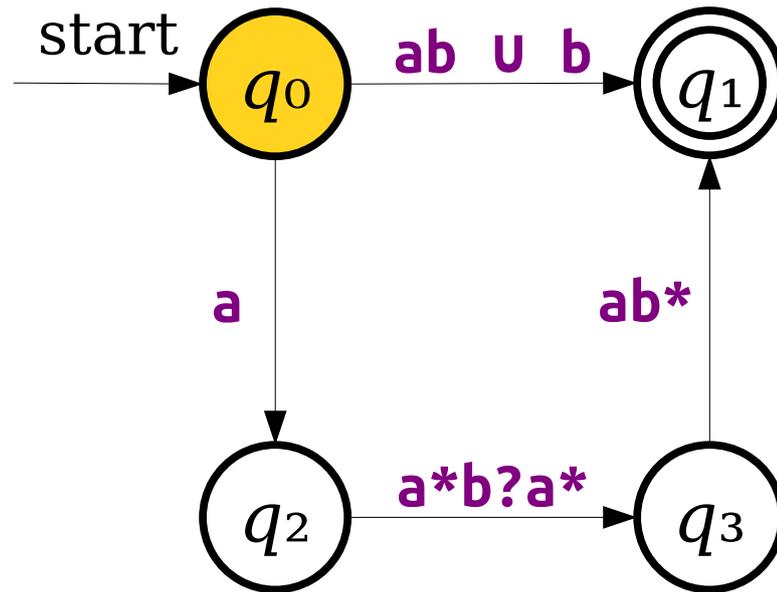
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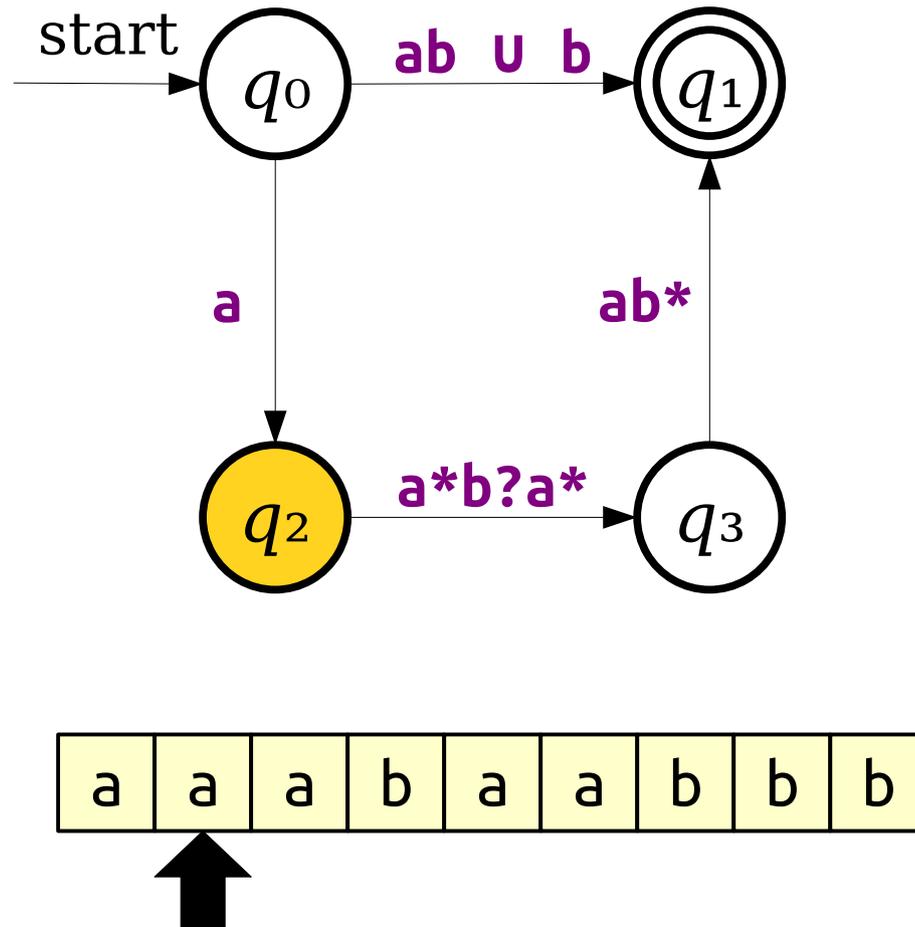
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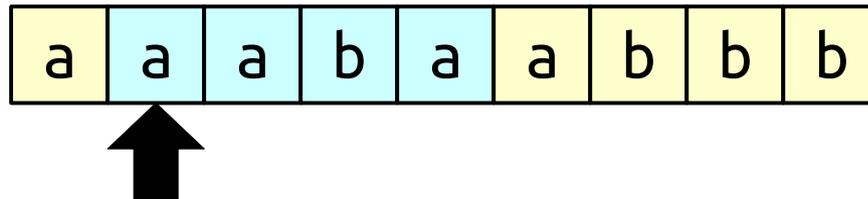
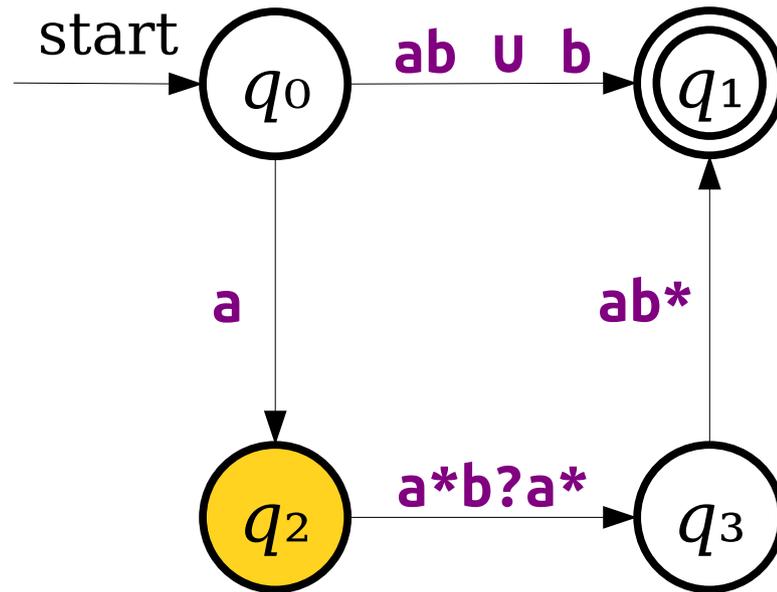
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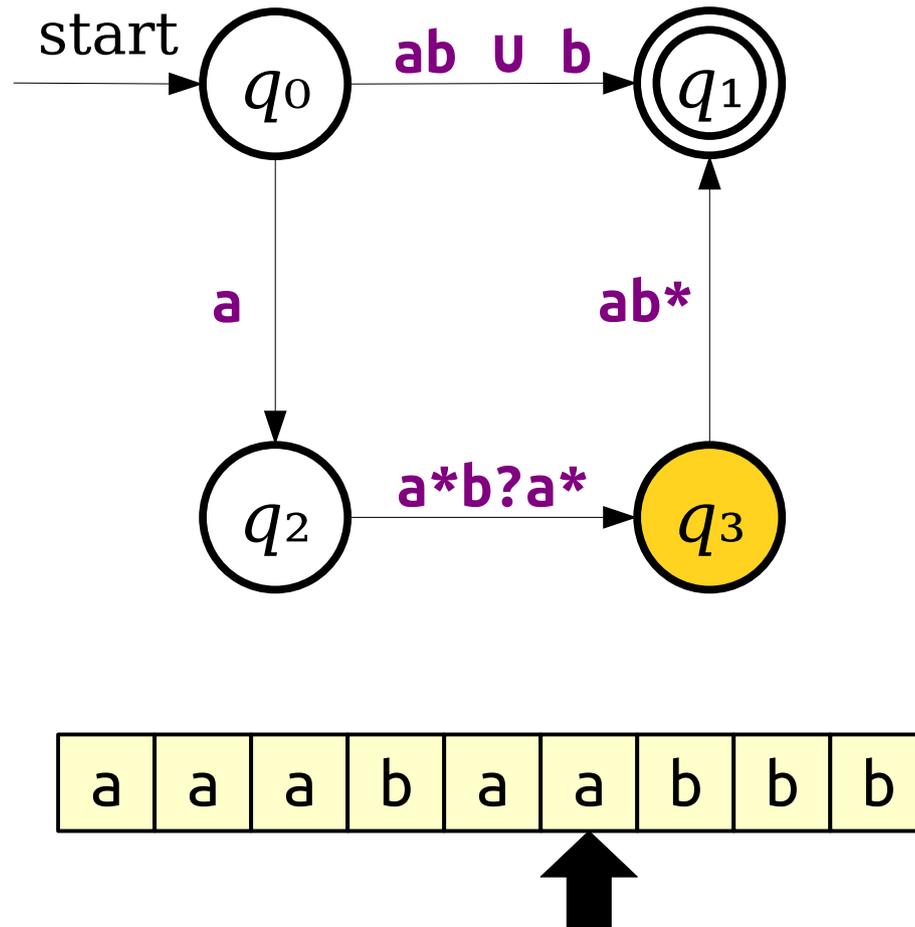
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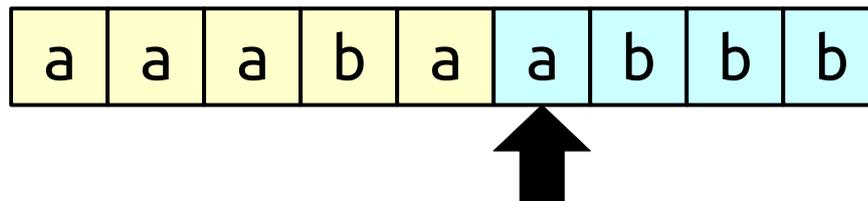
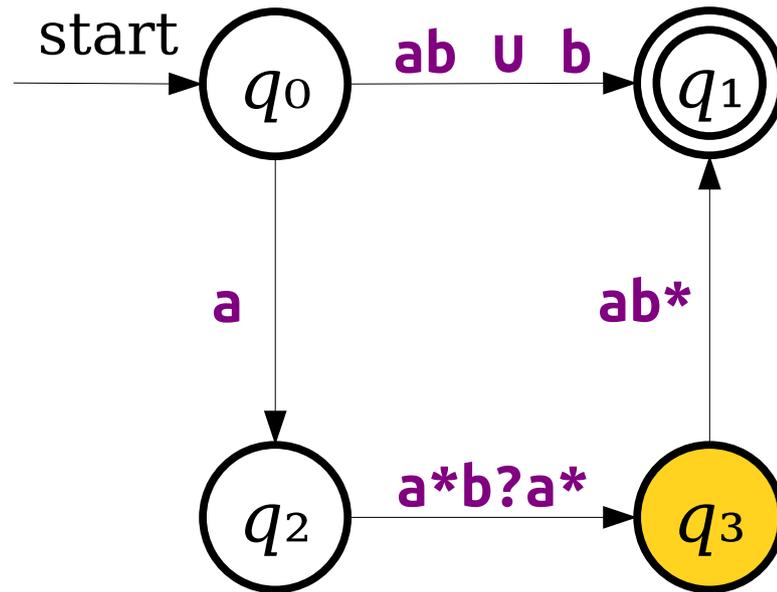
Generalizing NFAs



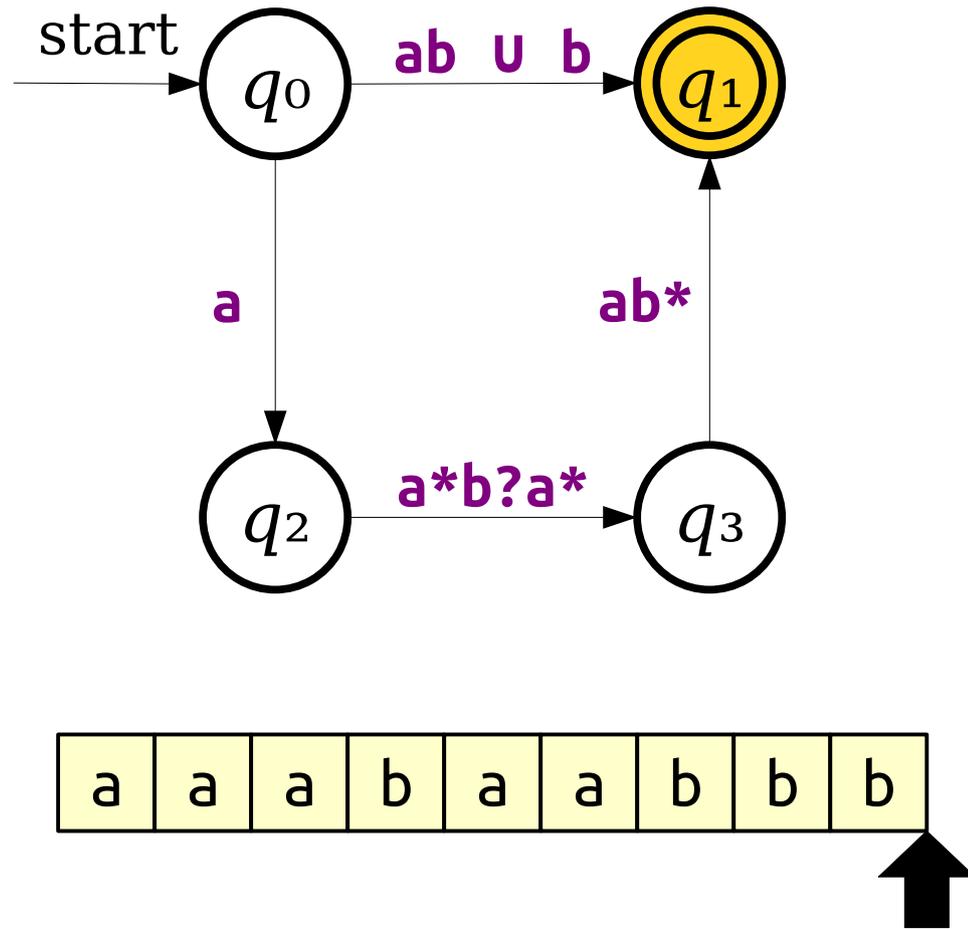
Generalizing NFAs



Generalizing NFAs



Generalizing NFAs



Key Idea 1: Imagine that we can label transitions in an NFA with arbitrary regular expressions.

Generalizing NFAs

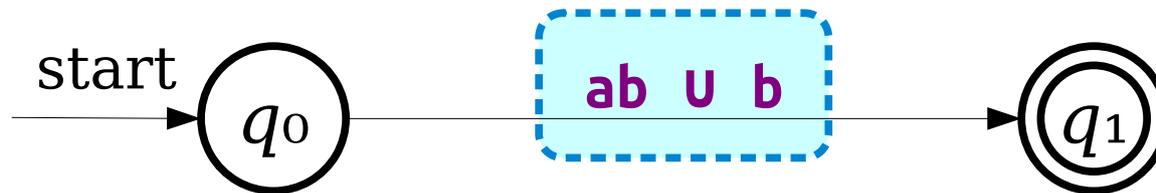


Generalizing NFAs



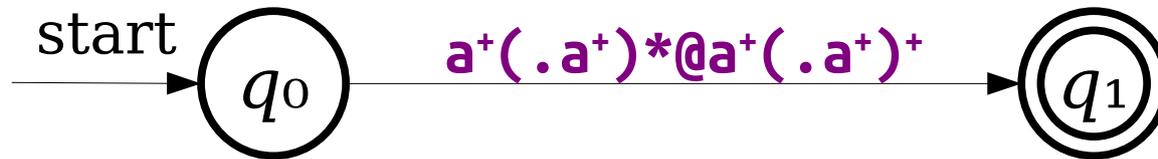
Is there a simple regular expression for the language of this generalized NFA?

Generalizing NFAs

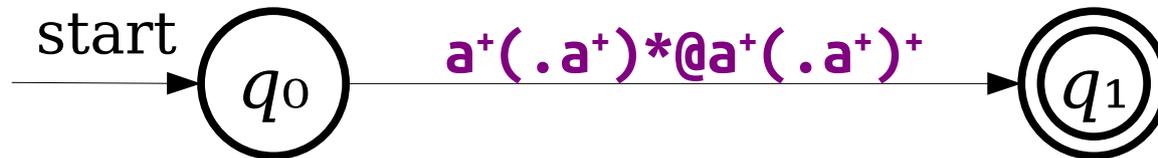


Is there a simple regular expression for the language of this generalized NFA?

Generalizing NFAs

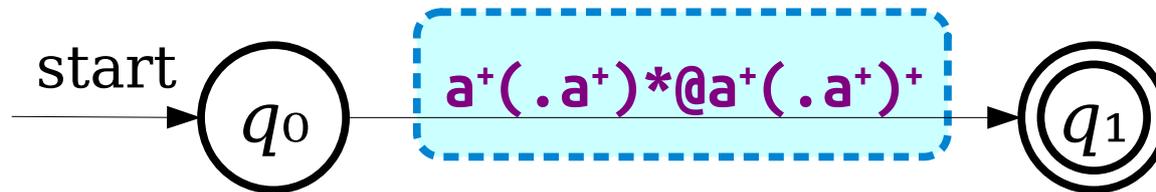


Generalizing NFAs



Is there a simple regular expression for the language of this generalized NFA?

Generalizing NFAs



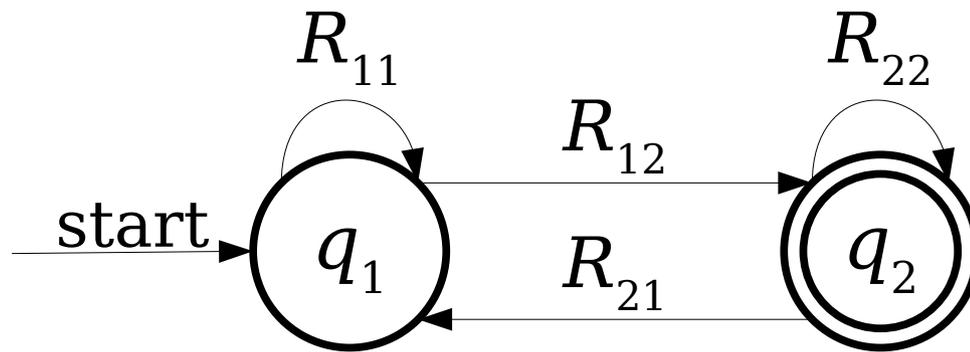
Is there a simple regular expression for the language of this generalized NFA?

Key Idea 2: If we can convert an NFA into a generalized NFA that looks like this...

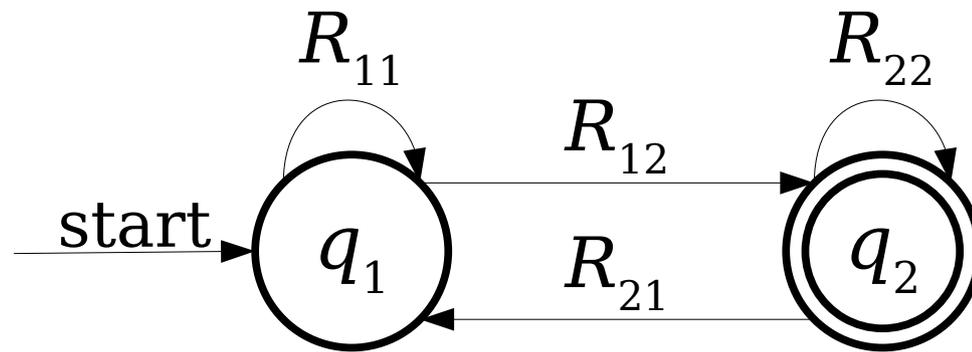


...then we can easily read off a regular expression for the original NFA.

From NFAs to Regular Expressions

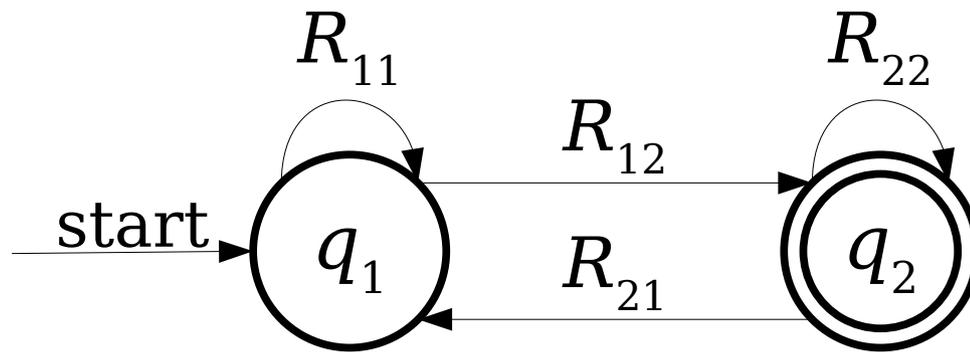


From NFAs to Regular Expressions



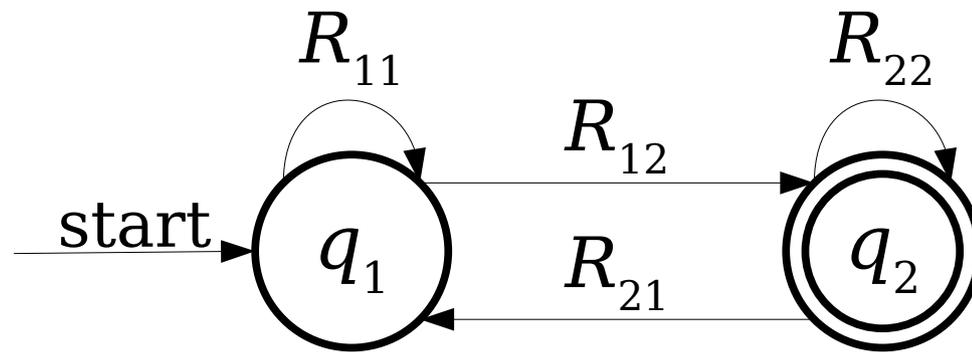
Here, R_{11} , R_{12} , R_{21} , and R_{22} are arbitrary regular expressions.

From NFAs to Regular Expressions



Question: Can we get a clean regular expression from this NFA?

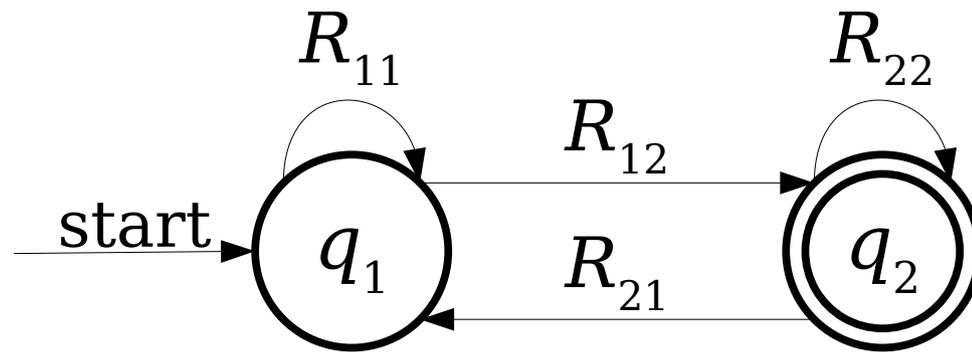
From NFAs to Regular Expressions



Key Idea 3: Somehow transform this NFA so that it looks like this:

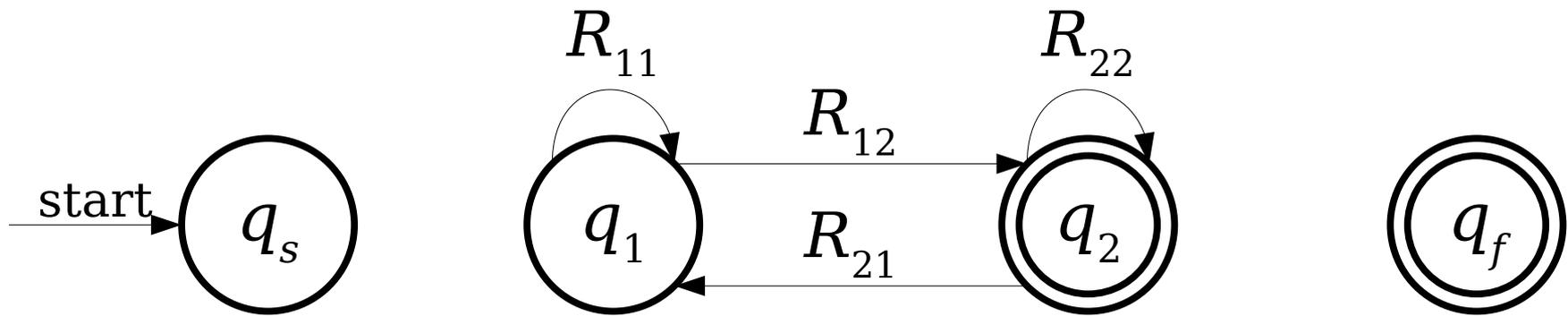


From NFAs to Regular Expressions

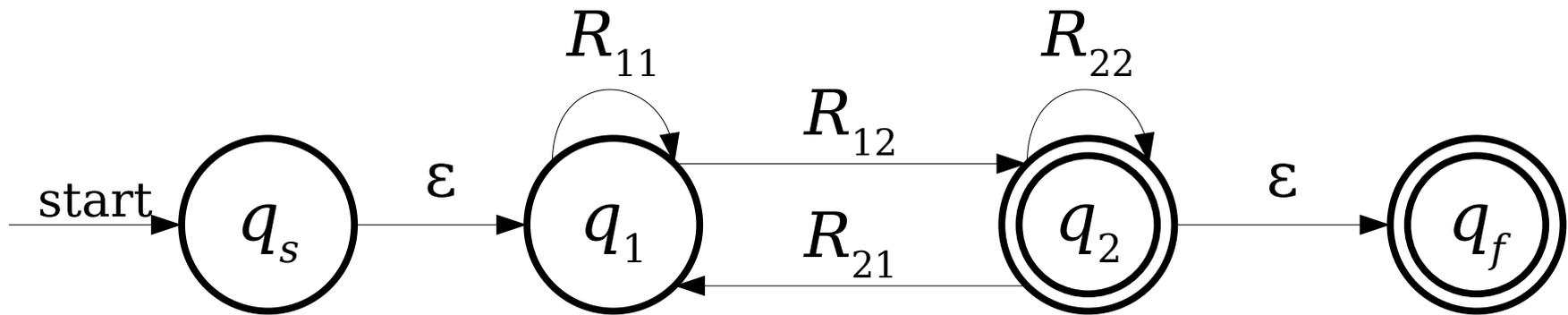


The first step is going to be a bit weird...

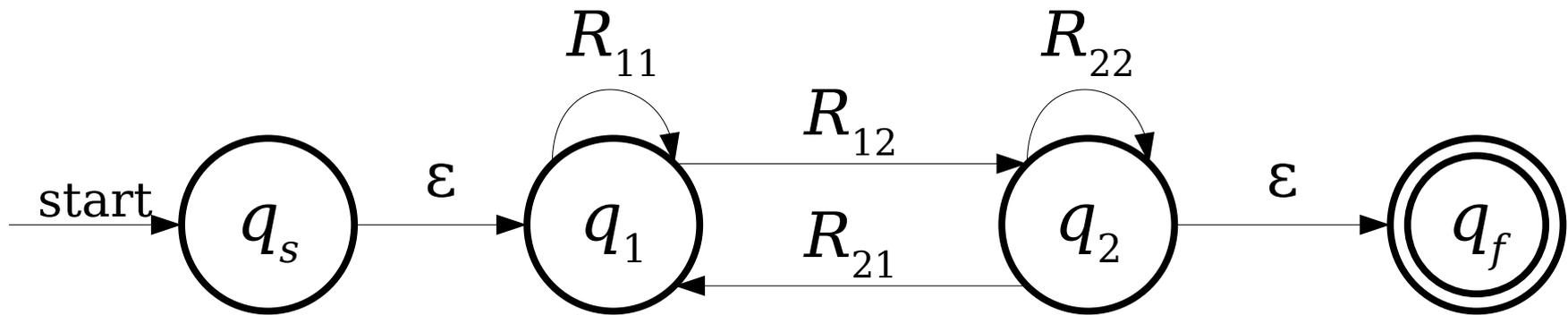
From NFAs to Regular Expressions



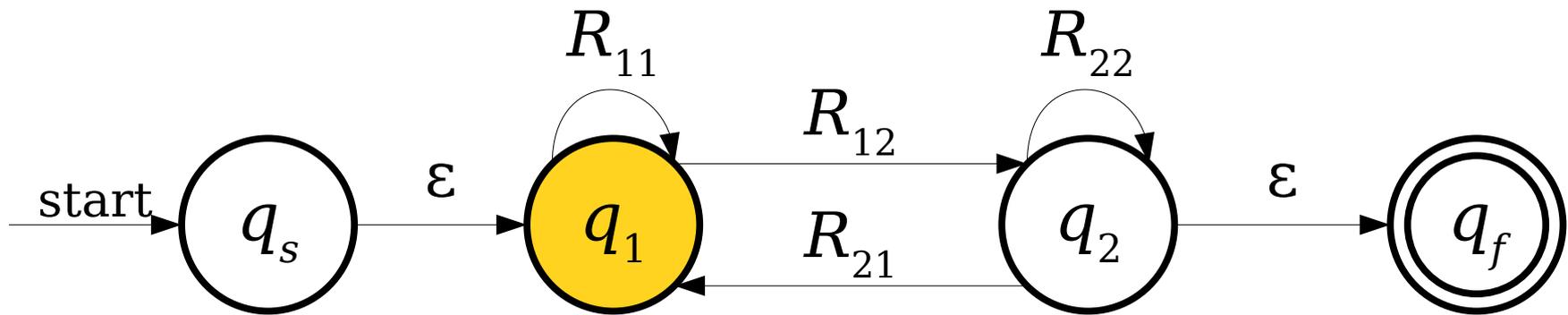
From NFAs to Regular Expressions



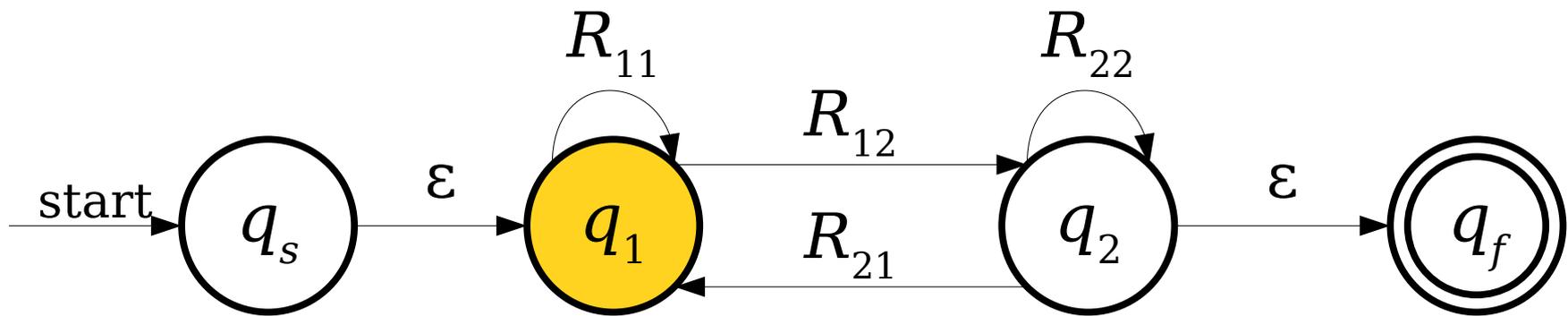
From NFAs to Regular Expressions



From NFAs to Regular Expressions

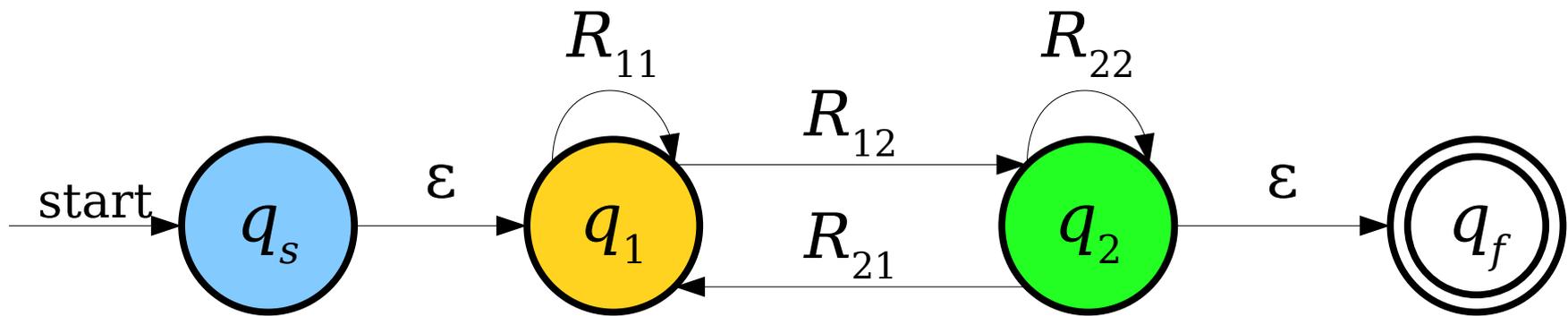


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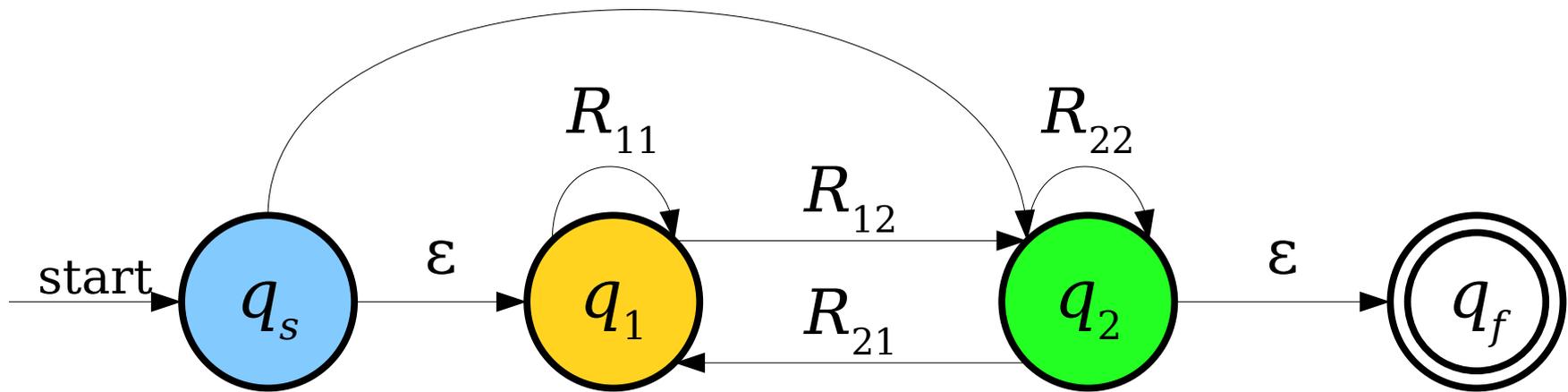


Could we eliminate
this state from
the NFA?

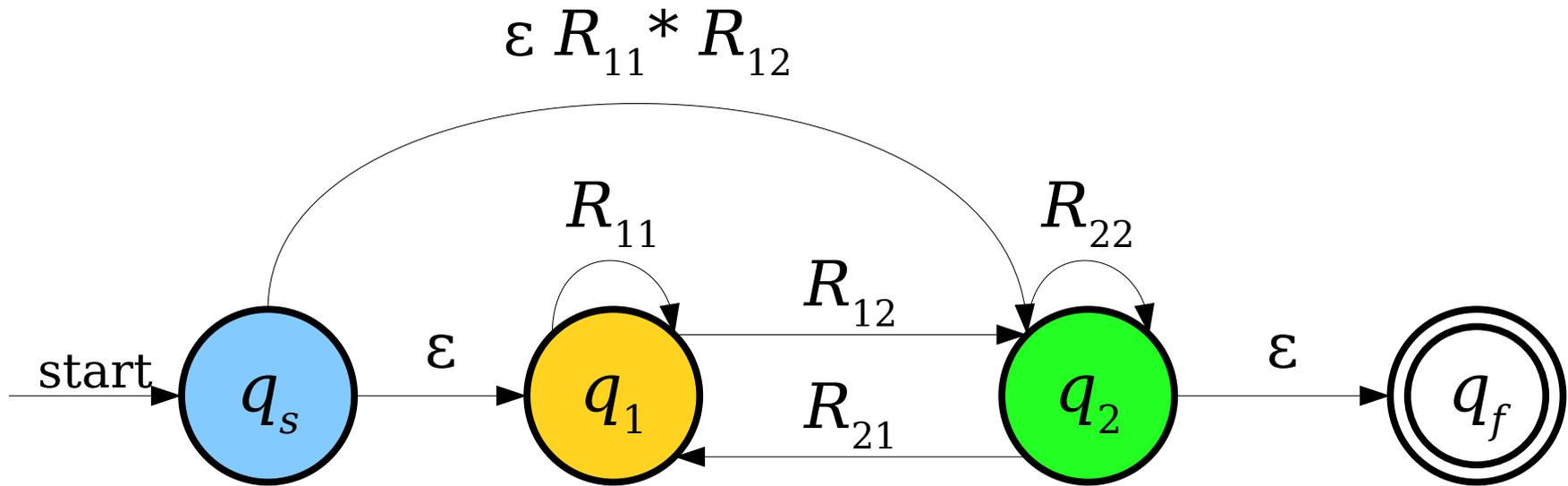
From NFAs to Regular Expressions



From NFAs to Regular Expressions

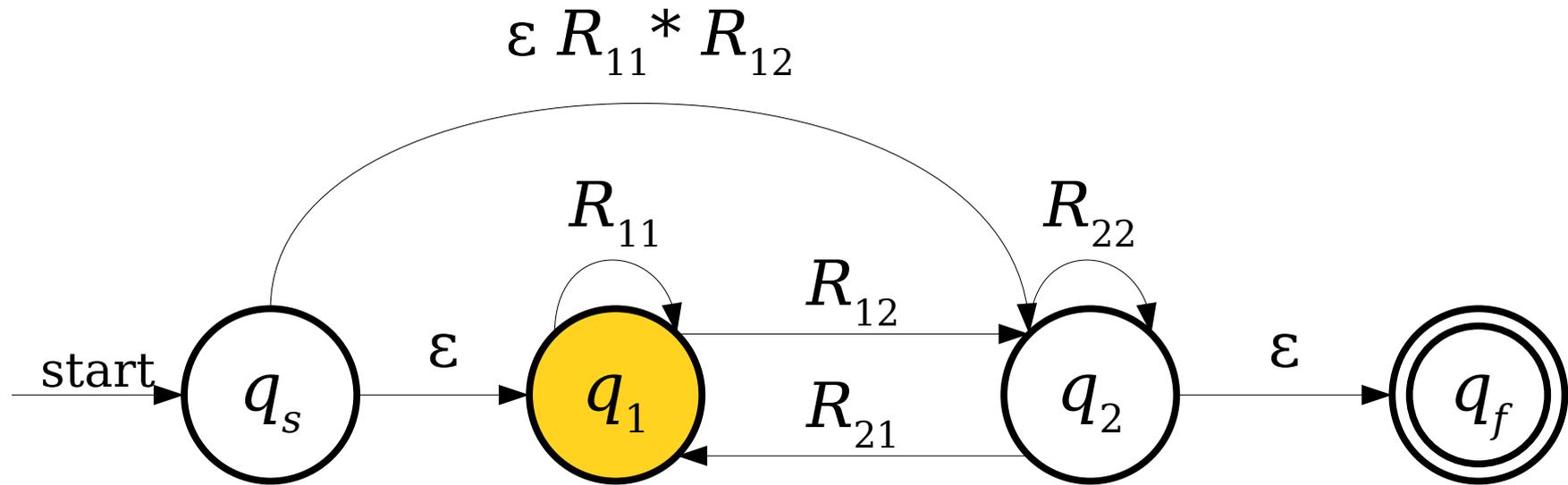


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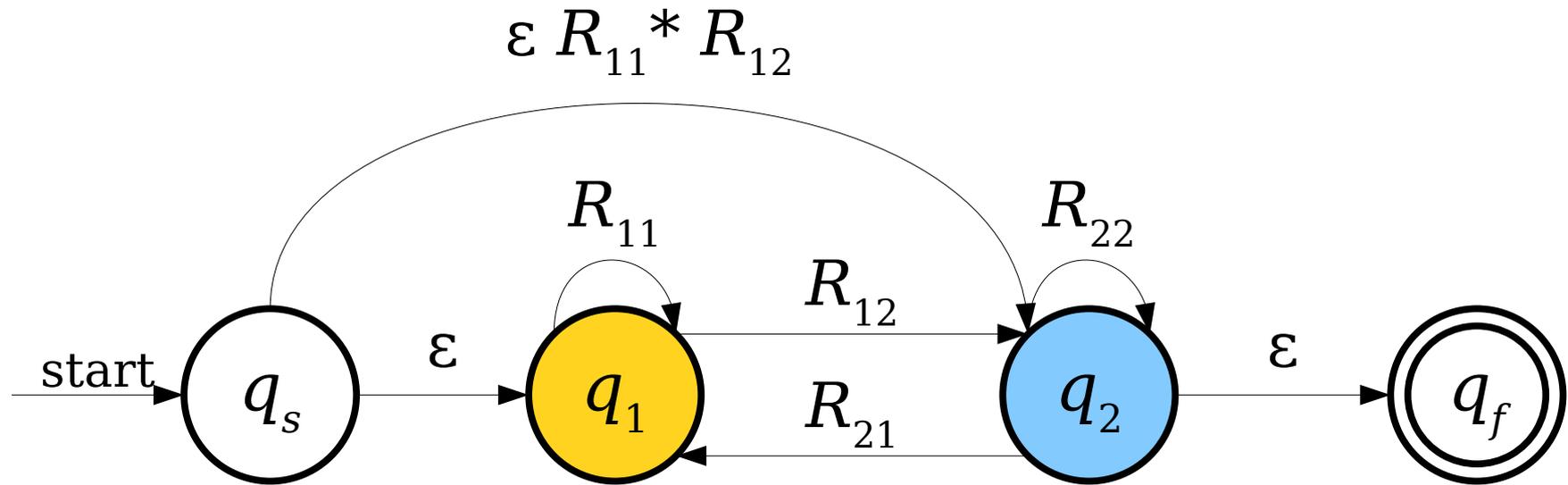


Note: We're using concatenation and Kleene closure in order to skip this state.

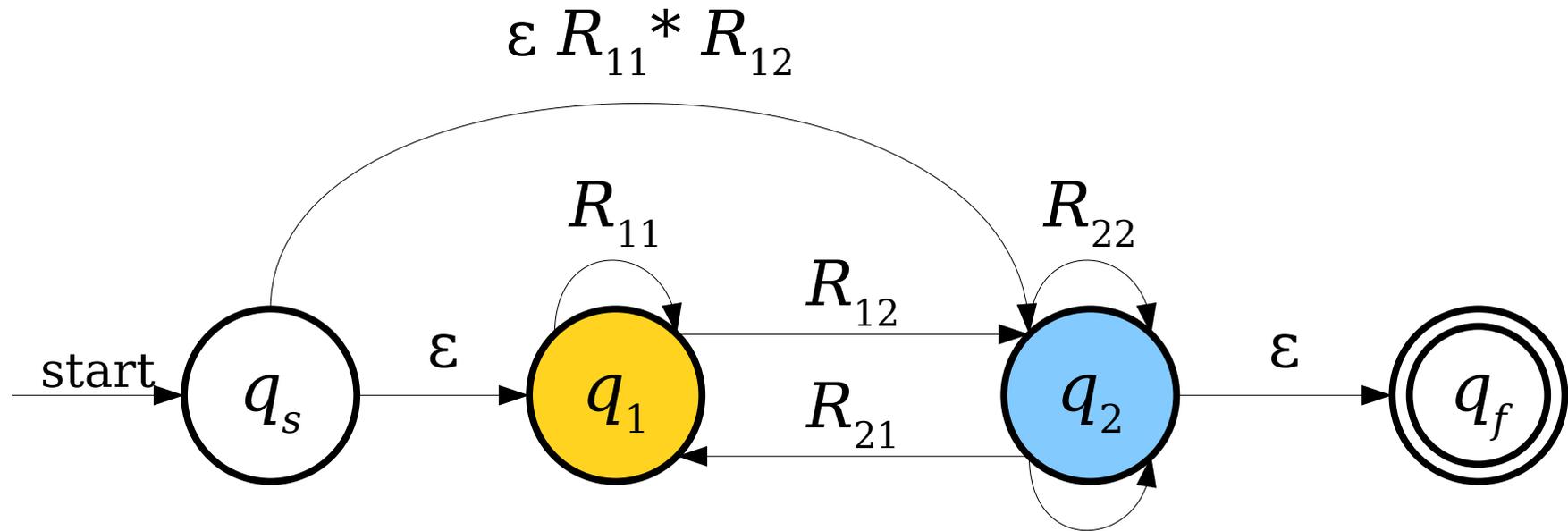
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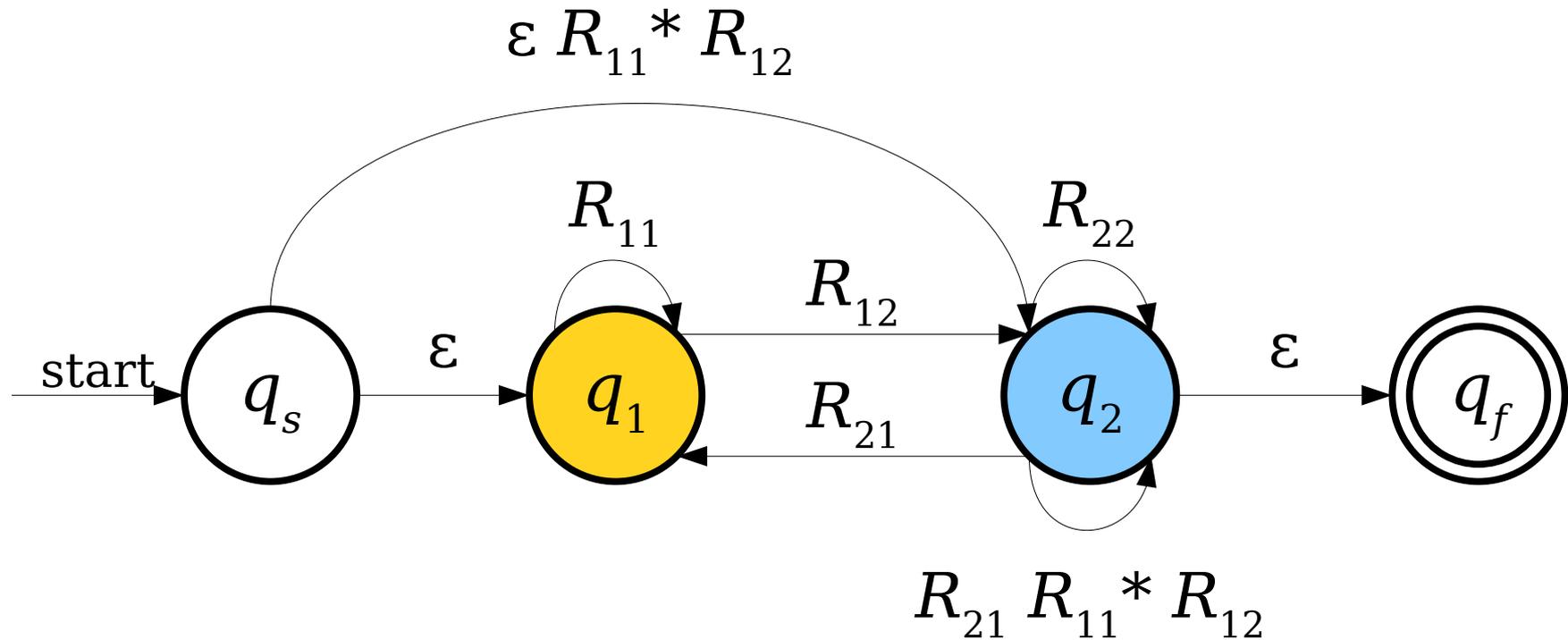
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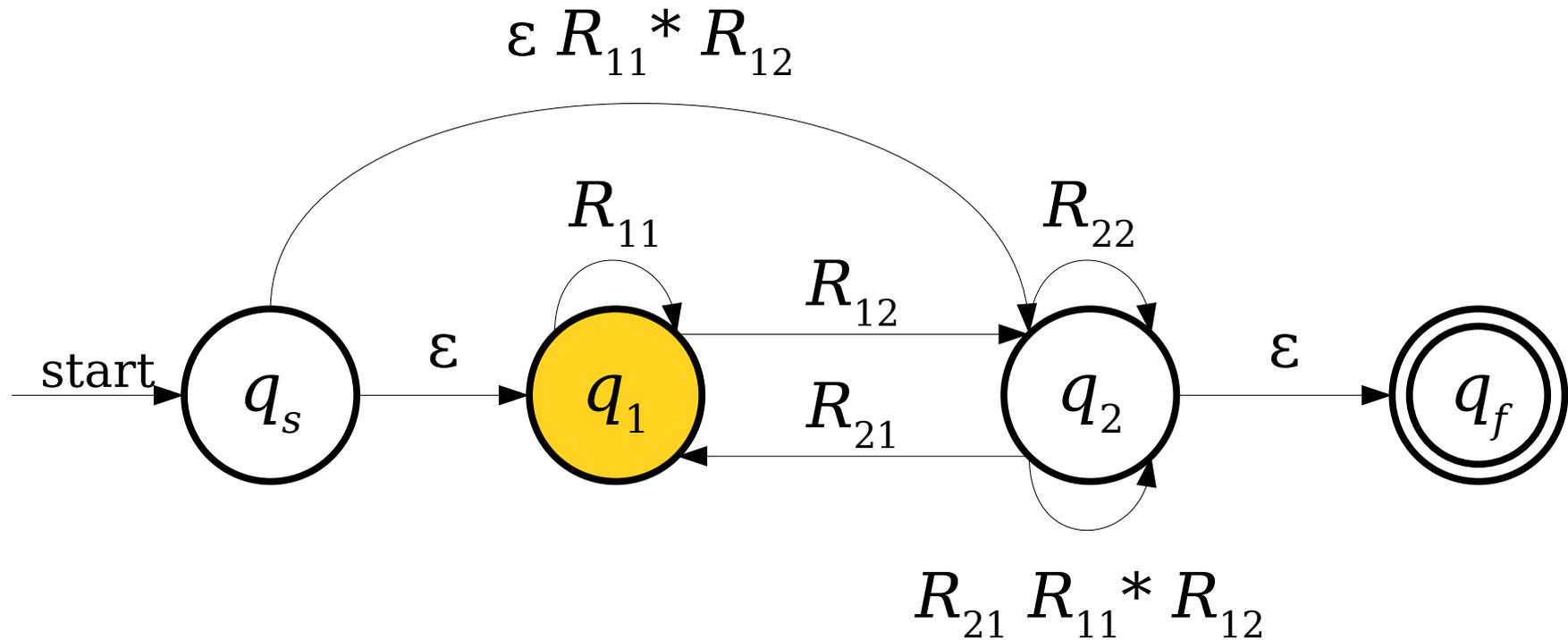
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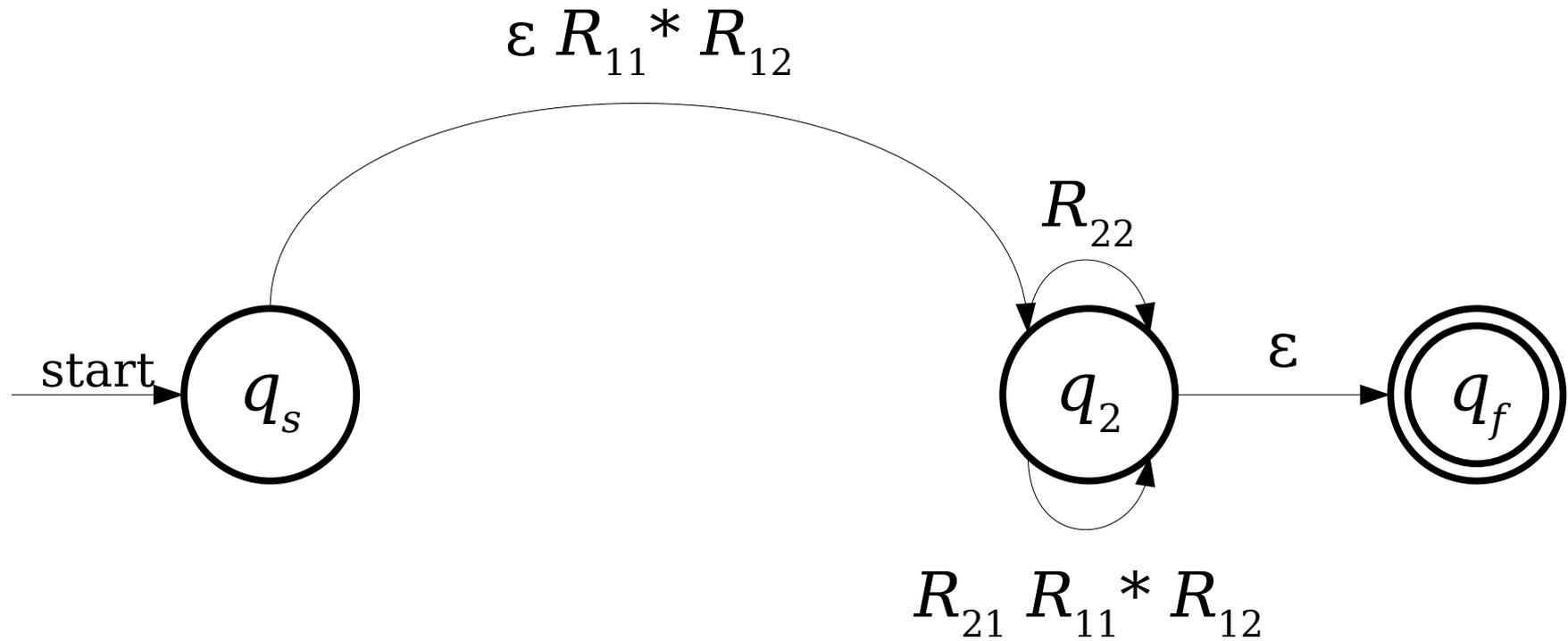
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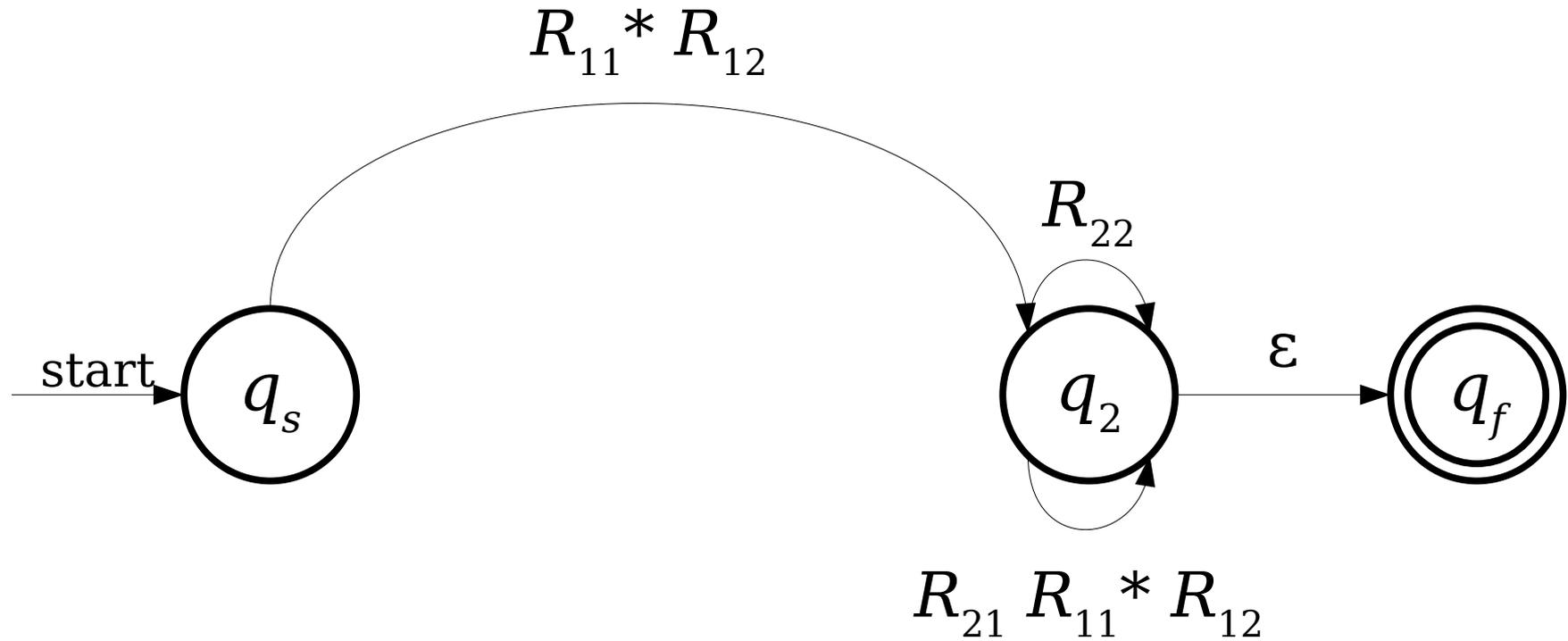
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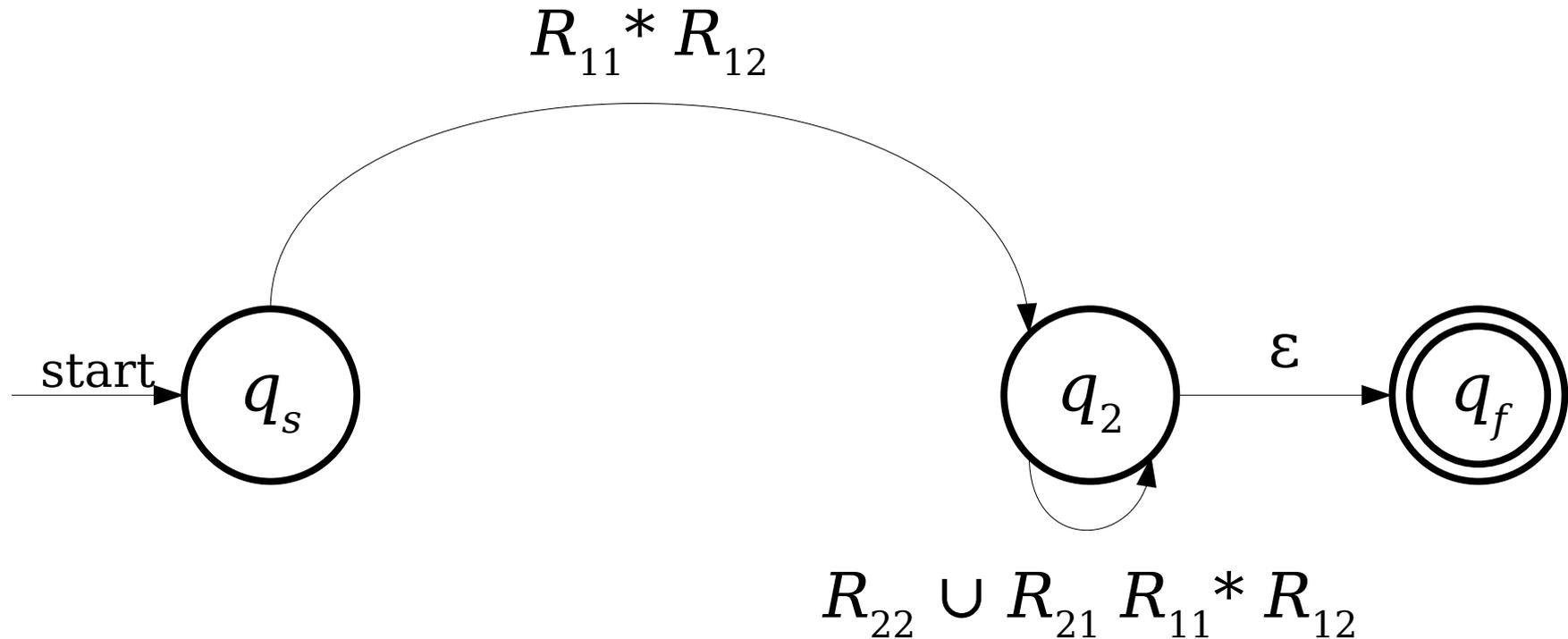
From NFAs to Regular Expressions



From NFAs to Regular Expressions

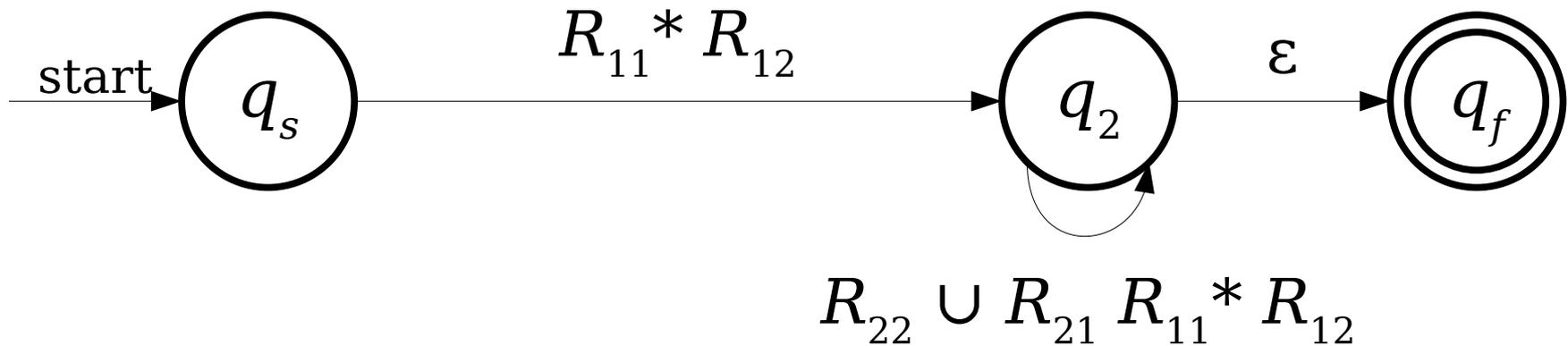


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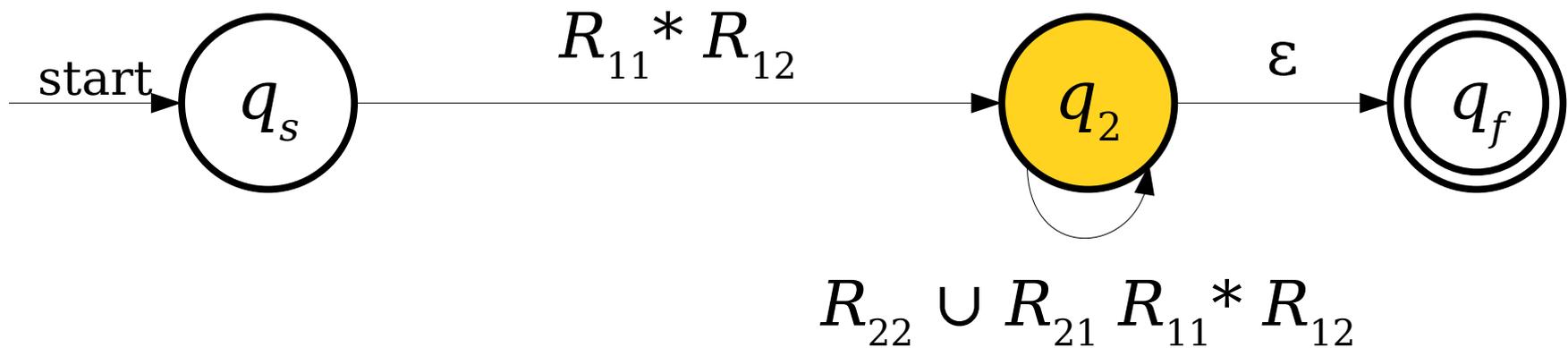


Note: We're using **union** to combine these transitions together.

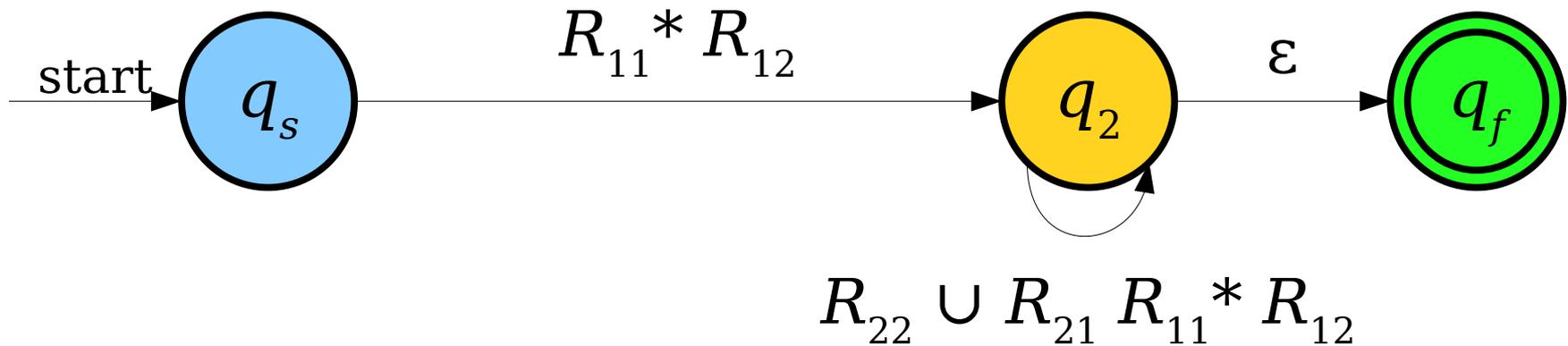
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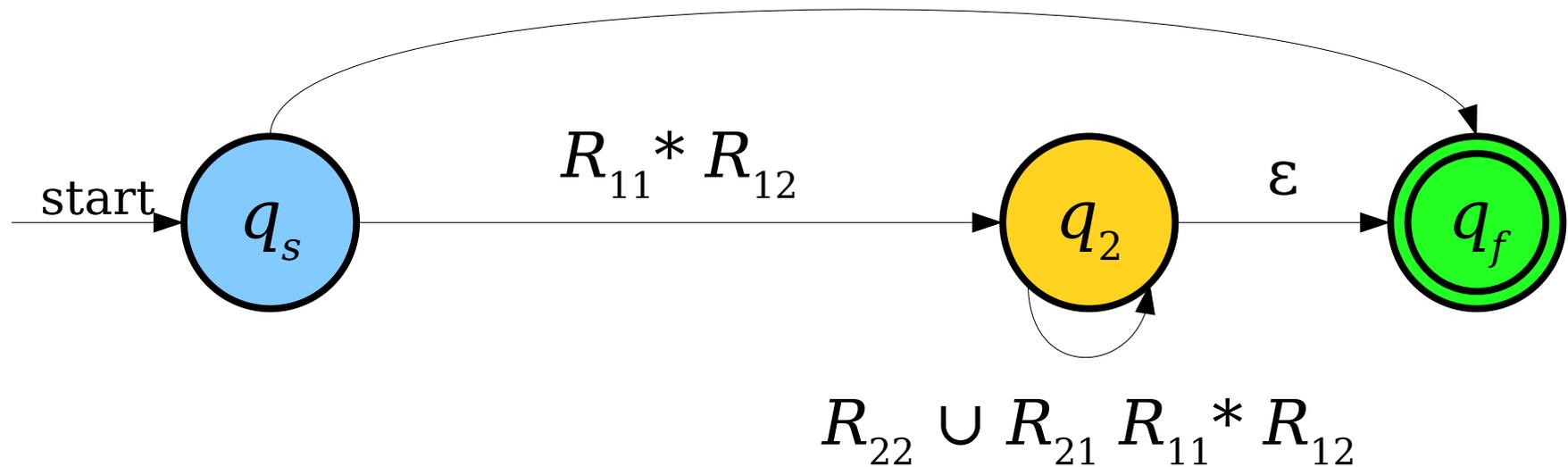
From NFAs to Regular Expressions



From NFAs to Regular Expressions

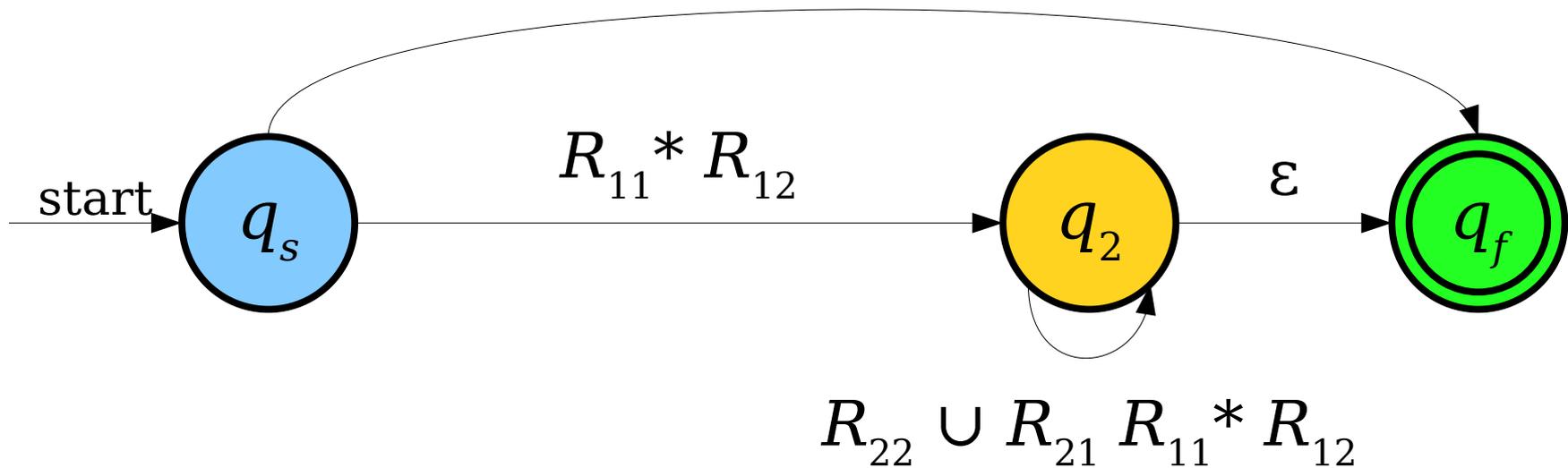


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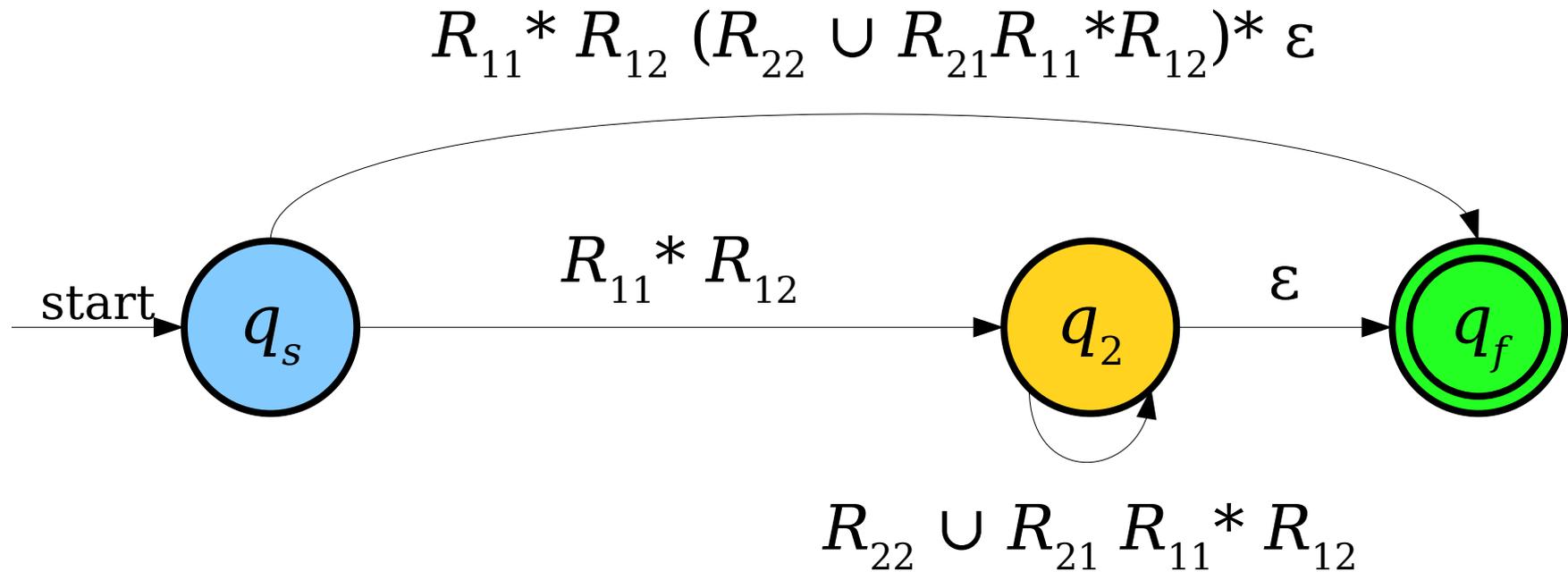


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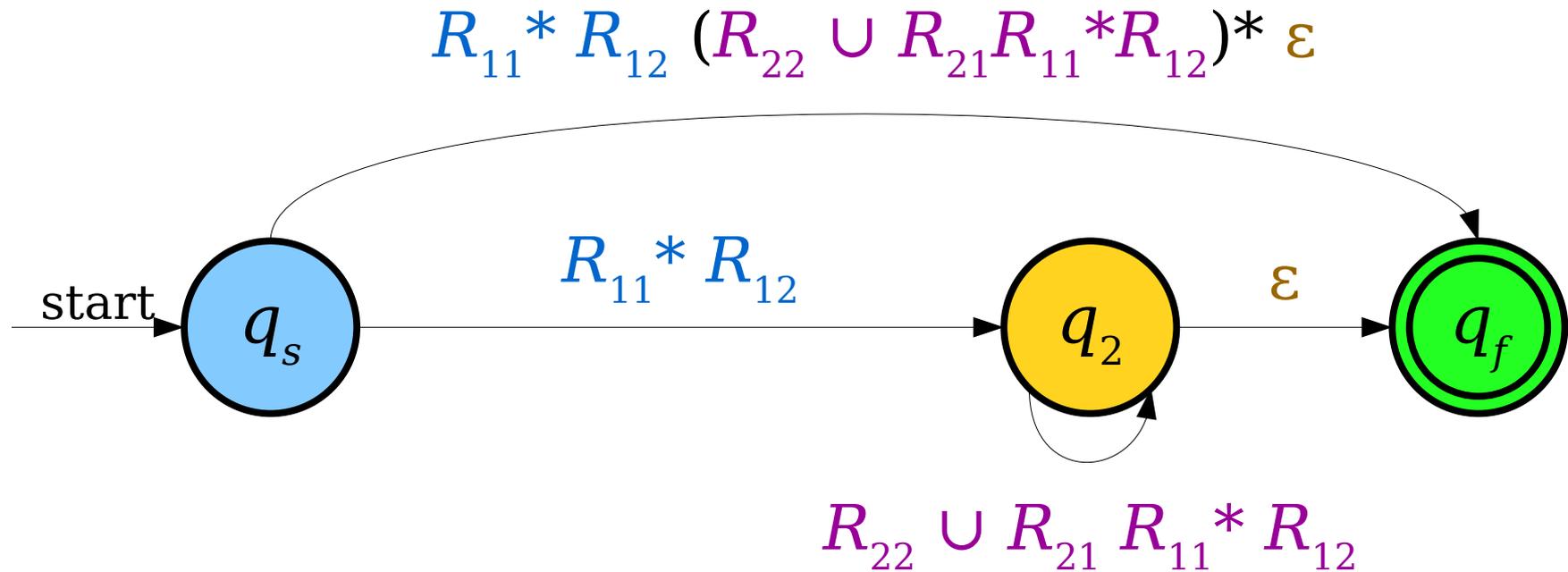
What should we put on this transition?



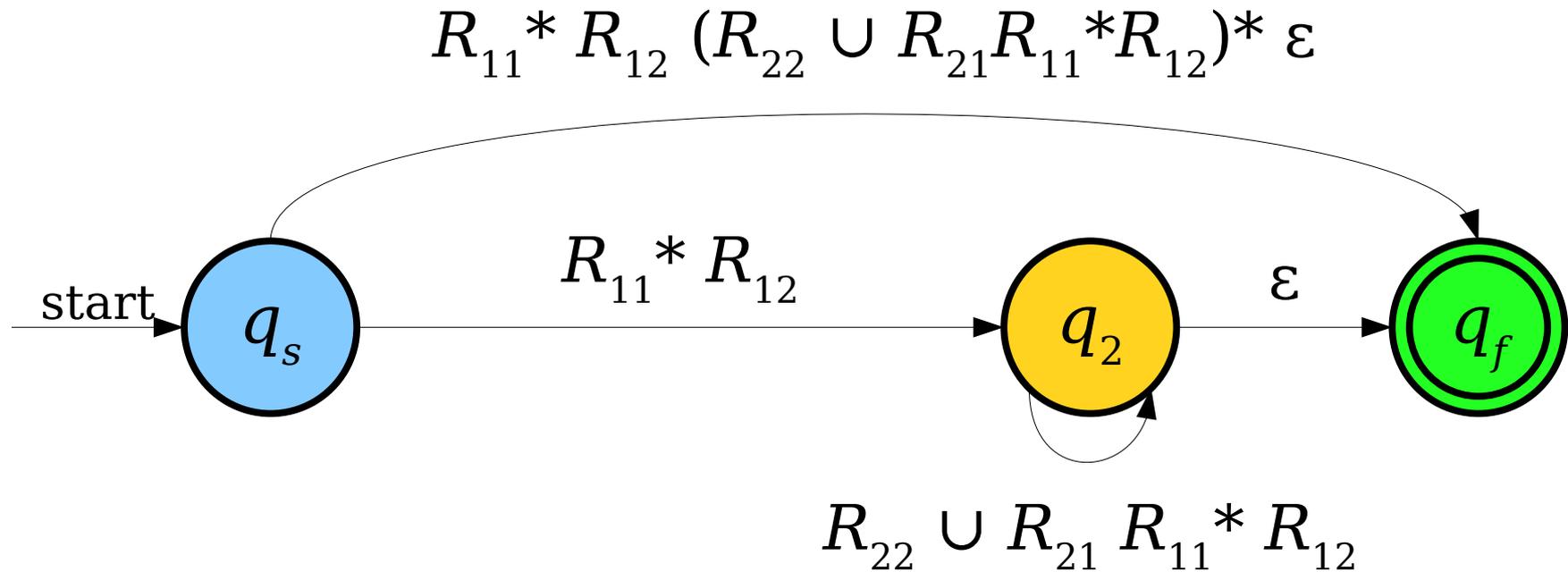
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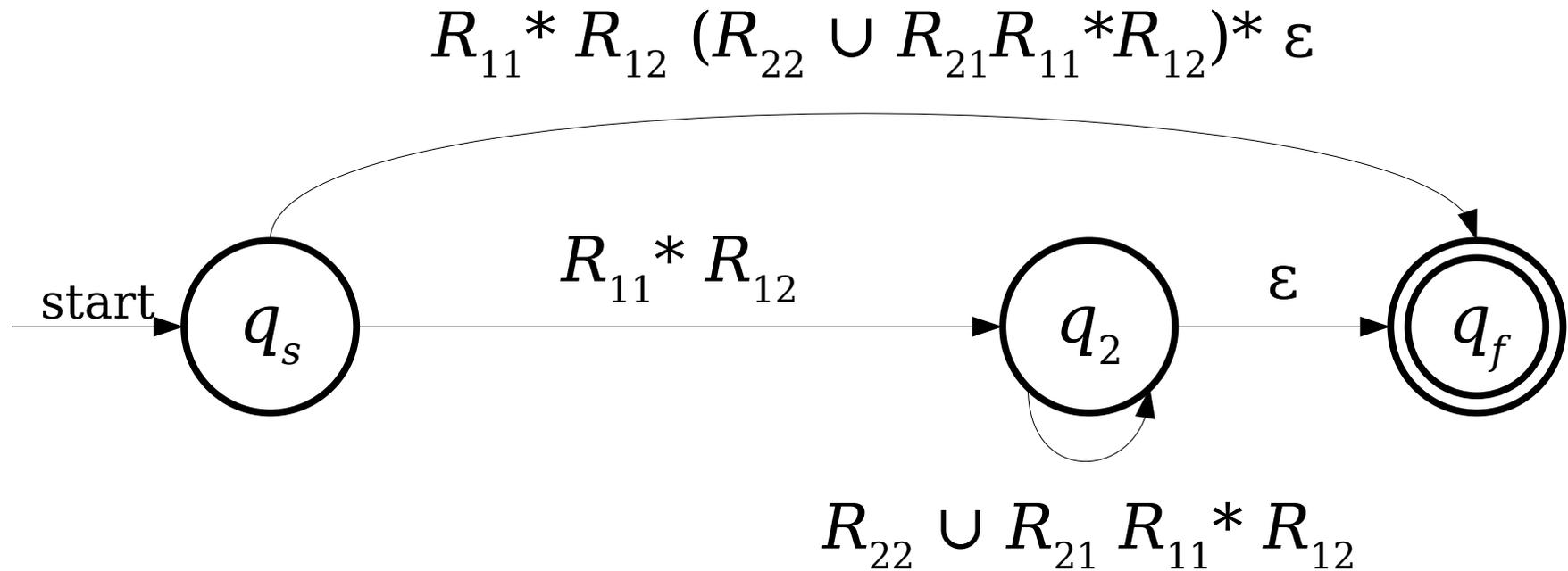
From NFAs to Regular Expressions



From NFAs to Regular Expressions

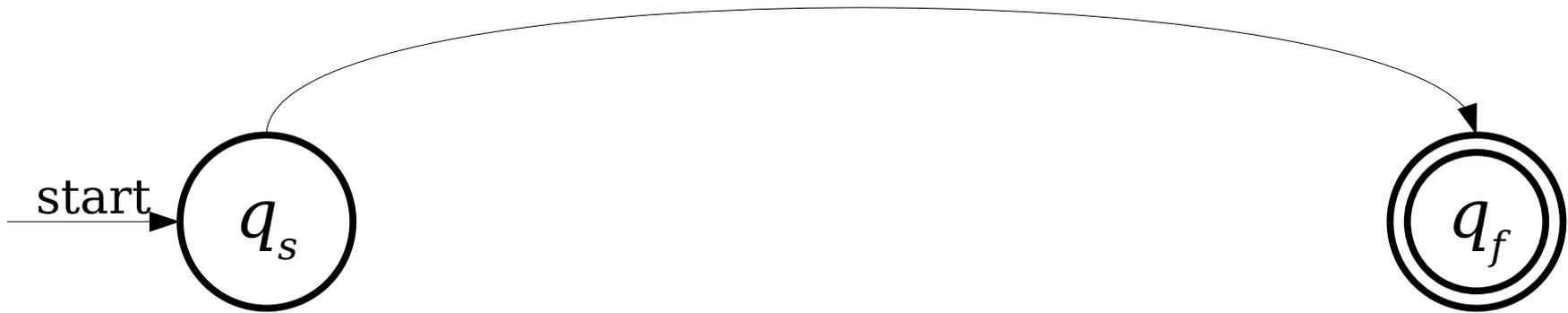


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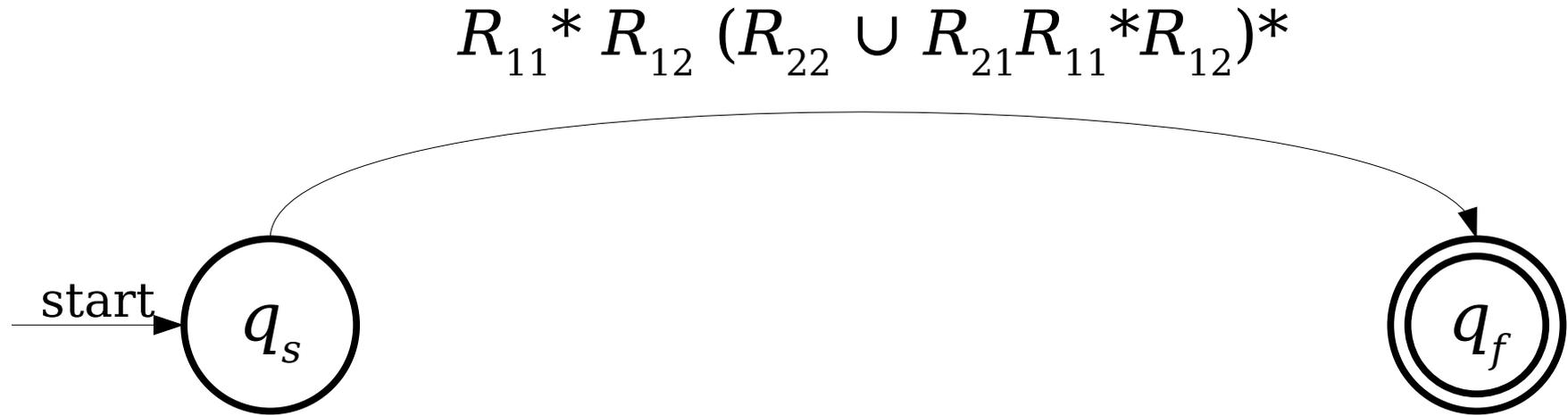


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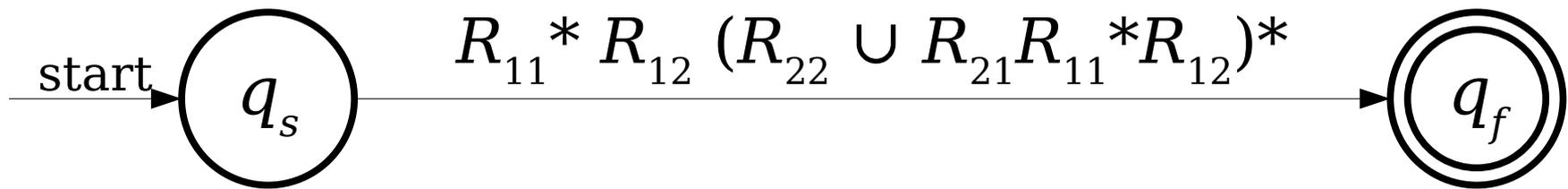
$$R_{11}^* R_{12} (R_{22} \cup R_{21} R_{11}^* R_{12})^* \varepsilon$$



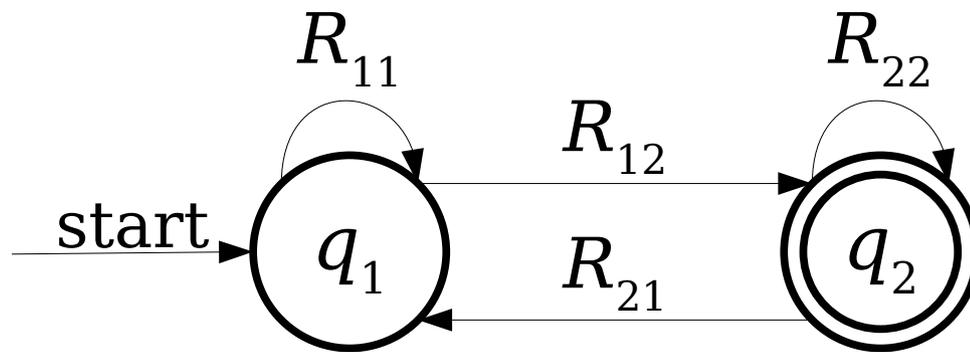
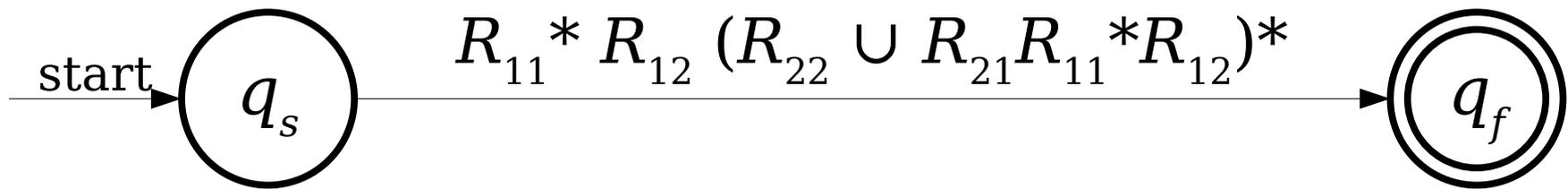
From NFAs to Regular Expressions



From NFAs to Regular Expressions



From NFAs to Regular Expressions



The State-Elimination Algorithm

- Start with an NFA N for the language L .
- Add a new start state q_s and accept state q_f to the NFA.
 - Add an ε -transition from q_s to the old start state of N .
 - Add ε -transitions from each accepting state of N to q_f , then mark them as not accepting.
- Repeatedly remove states other than q_s and q_f from the NFA by “shortcutting” them until only two states remain: q_s and q_f .
- The transition from q_s to q_f is then a regular expression for the NFA.

The State-Elimination Algorithm

- To eliminate a state q from the automaton, do the following for each pair of states q_0 and q_1 , where there's a transition from q_0 into q and a transition from q into q_1 :
 - Let R_{in} be the regex on the transition from q_0 to q .
 - Let R_{out} be the regex on the transition from q to q_1 .
 - If there is a regular expression R_{stay} on a transition from q to itself, add a new transition from q_0 to q_1 labeled $((R_{in})(R_{stay})^*(R_{out}))$.
 - If there isn't, add a new transition from q_0 to q_1 labeled $((R_{in})(R_{out}))$
- If a pair of states has multiple transitions between them labeled R_1, R_2, \dots, R_k , replace them with a single transition labeled $R_1 \cup R_2 \cup \dots \cup R_k$.